

# Language Technology for Internet Telephony Service Creation

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# Internet Telephony Services



- Rapidly evolving with addition of new functionality such email, database access and web services

## Service Creation



# Issues in Service Development

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- ◆ Software Intensive
  - Interaction with low-level hardware
  - Varying network types and capabilities
  - Customer Needs
- ◆ Robustness
- ◆ High Performance
  - Reuse Commodity services
  - Safety, security, dependability
- ◆ High Performance
  - Multimedia activities
  - Multiple processing layers

# Solution

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Programming Language  $\equiv$  Enabling Technology

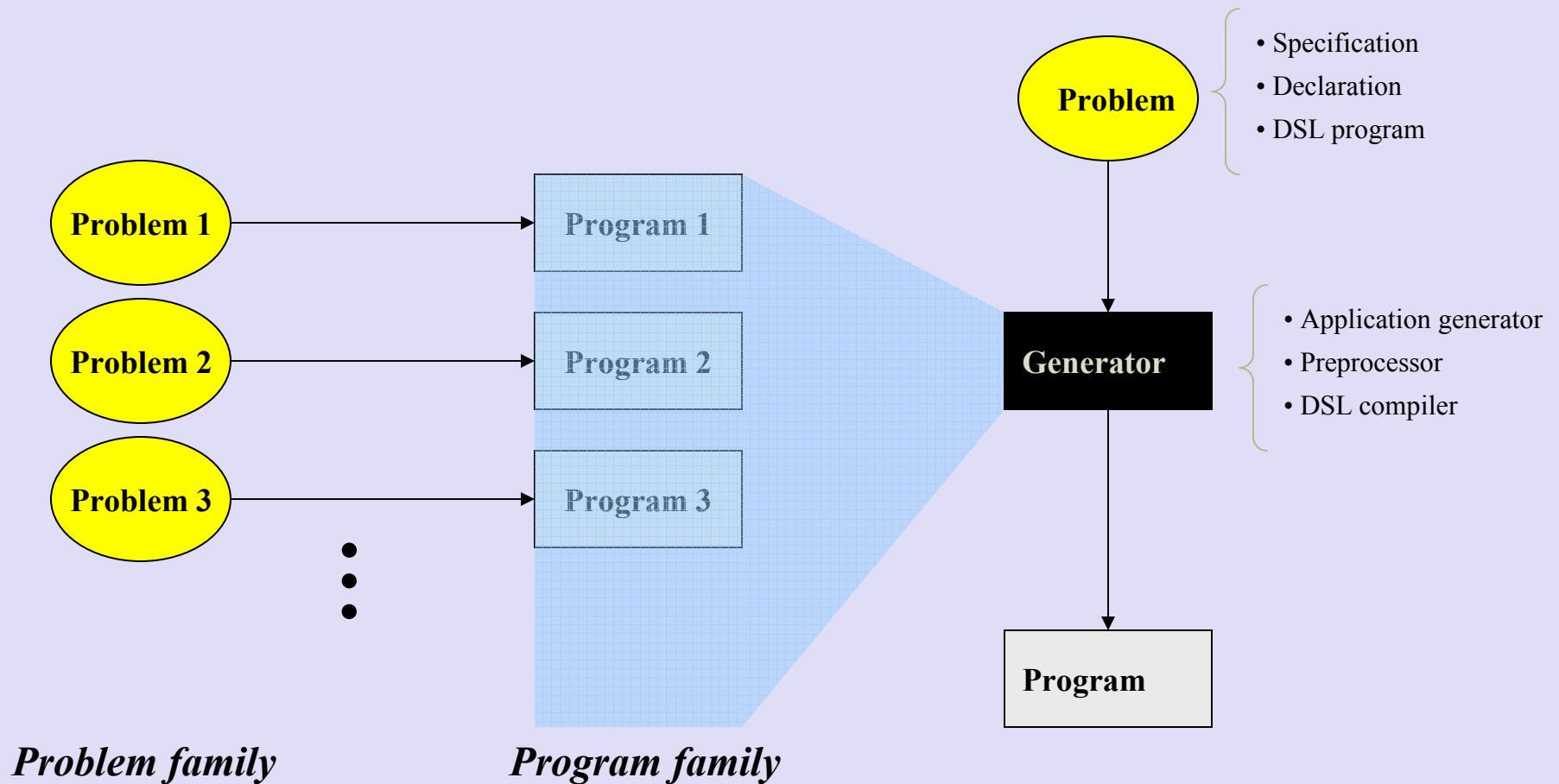
- ◆ Software intensive → Language design
- ◆ Robustness → Program analysis
- ◆ High Performance → Program transformation

# Talk Outline

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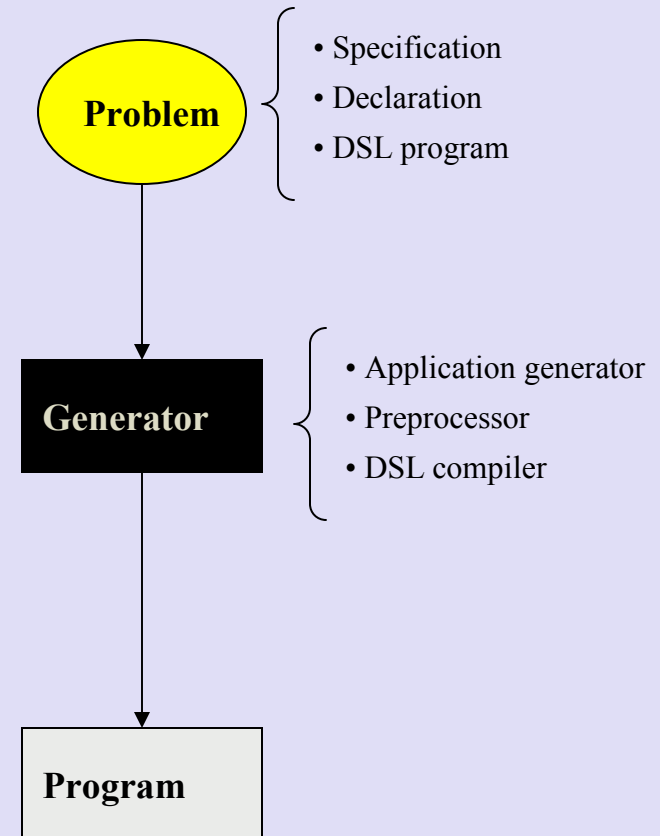
- ◆ Introduction to Domain Specific Languages (DSL)
- ◆ Overview of SIP
- ◆ SPL: A DSL for communication services
- ◆ Properties of SPL
- ◆ Summary

# Domain-Specific Languages: The Basic Idea



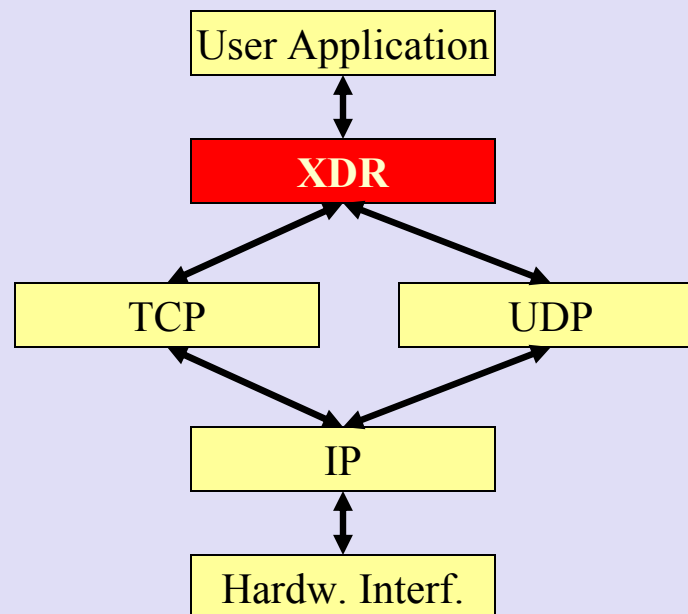
# Domain-Specific Languages: Issues

- ◆ When to develop a DSL?
- ◆ What is the scope of the DSL?
- ◆ How to design for the DSL?
  - **Key concept: program family**
- ◆ How to implement the DSL?
  - **Approach: stepwise methodology**
    - Data coding/marshalling components
    - Device interface layers
- ◆ How to assess the benefits?
  - **Key concept: program family**



# A Program Family with a DSL: Machine Independent Data Coding

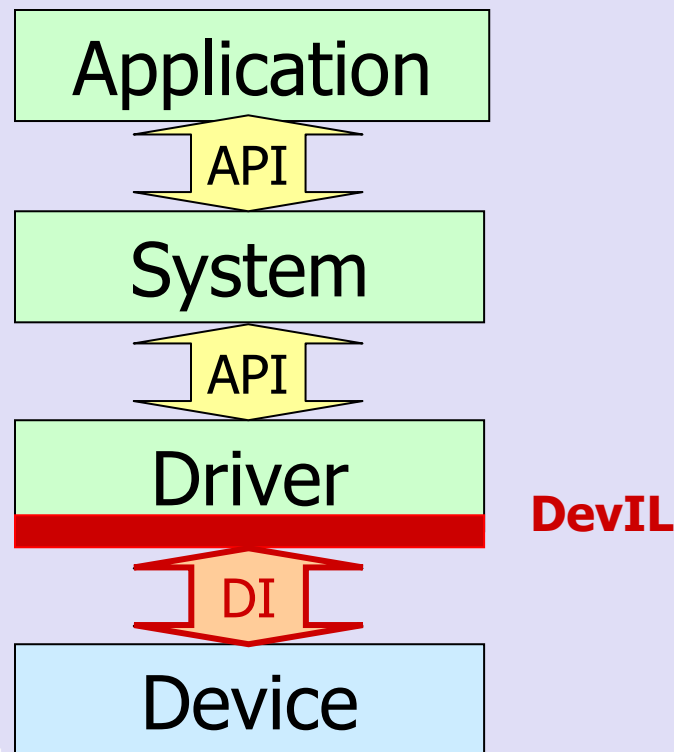
- ◆ Commonalities: data traversal, type description...
- ◆ Variations: coding direction, size, format...





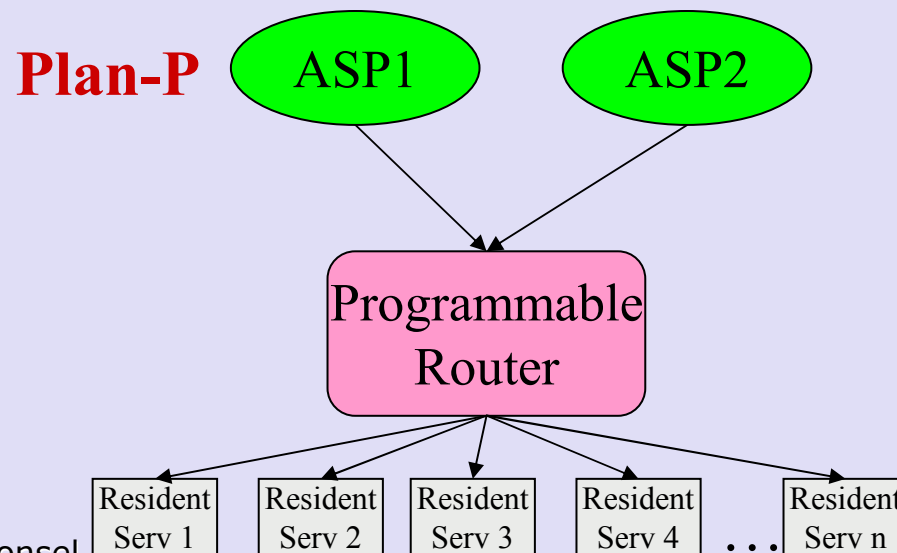
# A Program Family with a DSL: Device Driver Interface

- ◆ Commonalities: API, bit operations...
- ◆ Variations: parameters, registers...



# A Program Family with a DSL: Application-Specific Protocols

- ◆ Commonalities: API, packets operations
- ◆ Variations: routing policies, packet processing...



# Program Family Characterization: A Key Step Towards DSL Design

- ◆ Nature of program family
  - Existing (Devil)
  - Virtual ( $\approx$  Plan-P)
  - Both ( $\approx$  Bossa)
- ◆ Scope of programs
  - Family of layers (Devil)
  - Family of components (Yacc)
  - Family of systems (Scripting language)
- ◆ Range of commonalities  
(size of program family captured)
- ◆ Scope of computations  
(kind of expressible computations)

# Beyond Program Families

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- ◆ Technical literature
- ◆ Documentation
- ◆ Domain experts
- ◆ Current and future requirements

# Domain-Specific Languages: A **(Conceptual)** Definition

- ◆ Program family
- ◆ Domain-specific abstractions and notations
  - Conciseness, readability
- ◆ Declarative (often)
  - What to compute, not how to compute it
- ◆ Restricted/enriched semantics
  - Making critical properties decidable

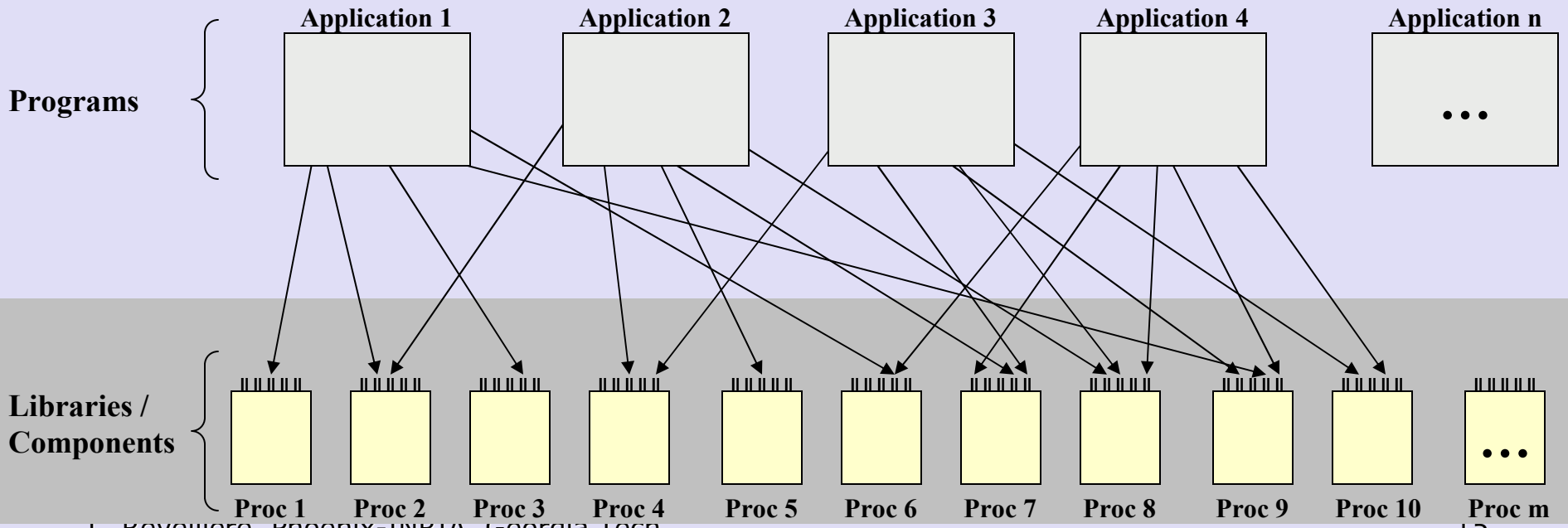
# Approach to DSLs

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- ◆ Resulting structure
  - Static semantics: interpreter/compiler
  - Dynamic semantics: Abstract machine
- ◆ Stepwise methodology:
  - Program family
  - Library/components
  - Abstract machine
  - DSL

# Program Family: Libraries / Components

- ◆ Commonalities  $\Rightarrow$  Library entries
- ◆ Variations  $\Rightarrow$  Parameters
- Bloated code
  - Numerous cases



# Coding Programs: XDR Library

`xdr_int`

`xdr_long`

`XDR_PUTLONG`

`xdrmem_putlong`

`htonl`

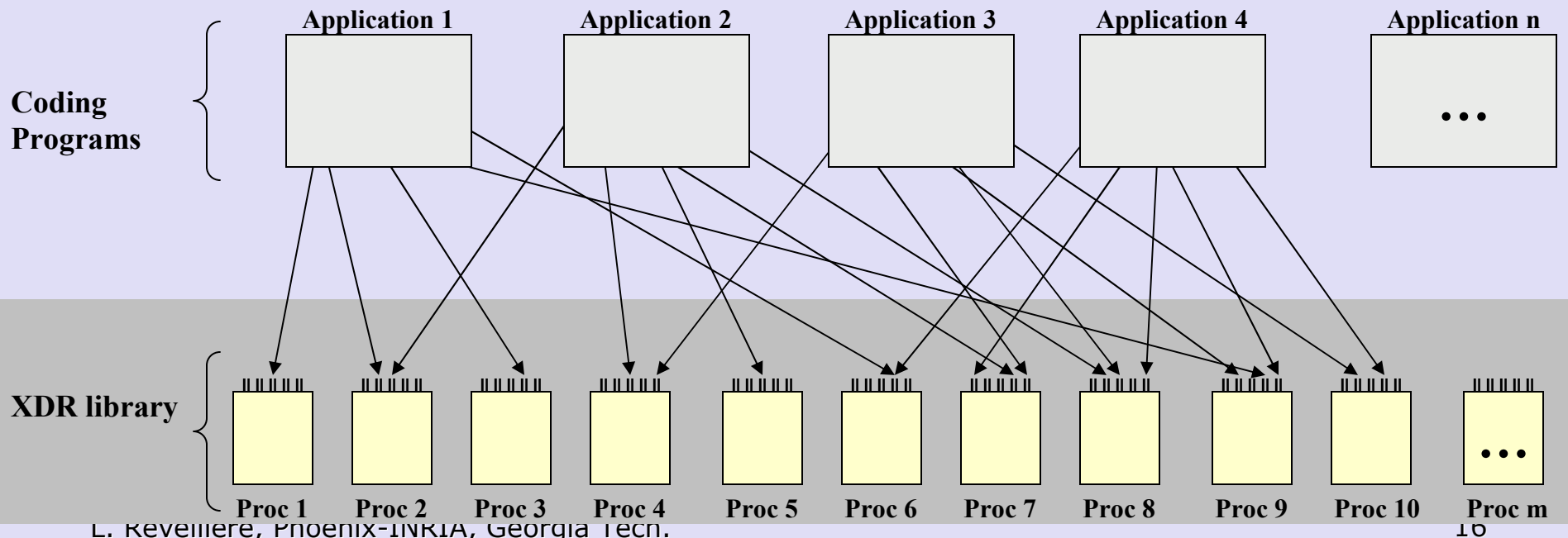
Integer size

Coding direction

Generic marshalling

Buffer write + overflow check

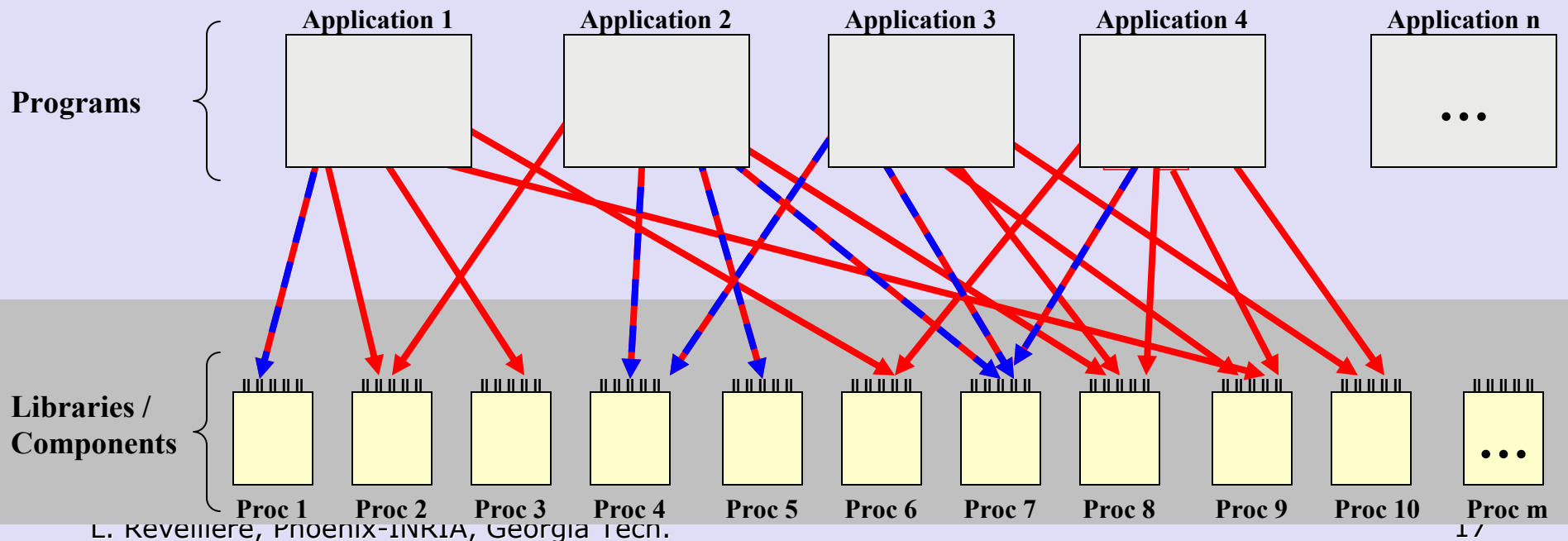
Byte ordering





# Program Family: Libraries / Components

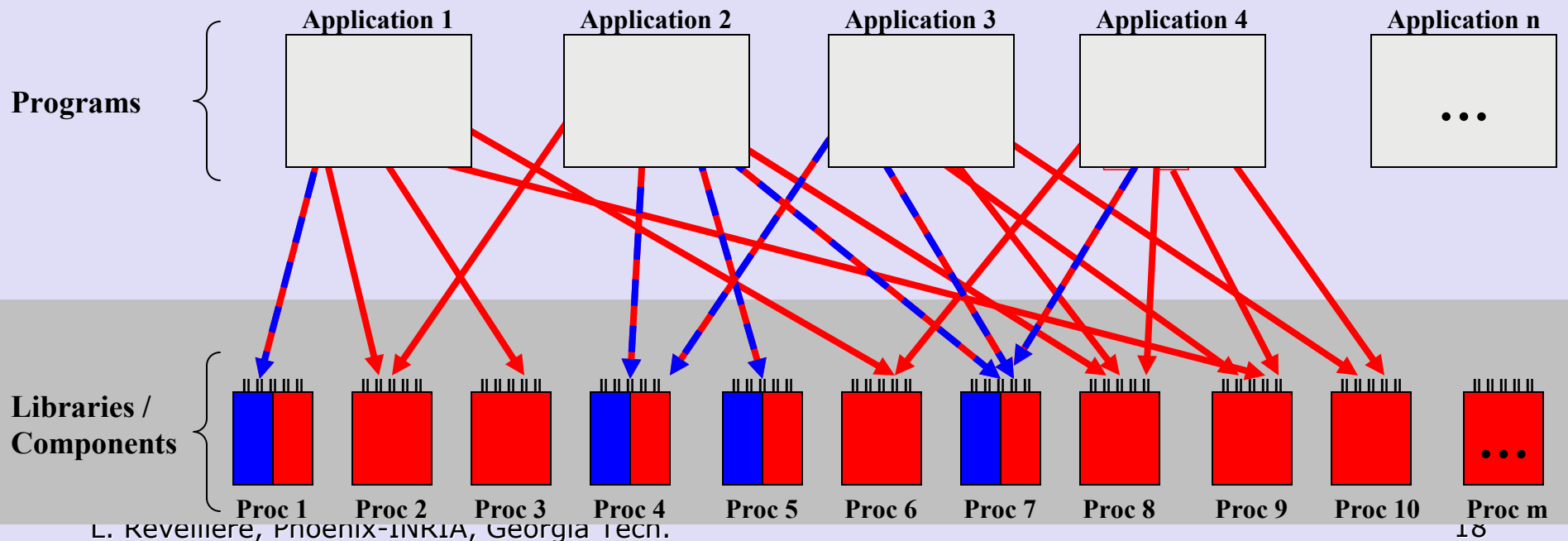
- For a given set of problems :  
*invariants*



# Libraries/Components :

## *Need For Customization*

- For a given set of problems :
  - Optimize time (and/or)
  - Optimize space



# XDR Library :

## *Need For Customization*

`xdr_int`

Integer size

`xdr_long`

Coding direction

`XDR_PUTLONG`

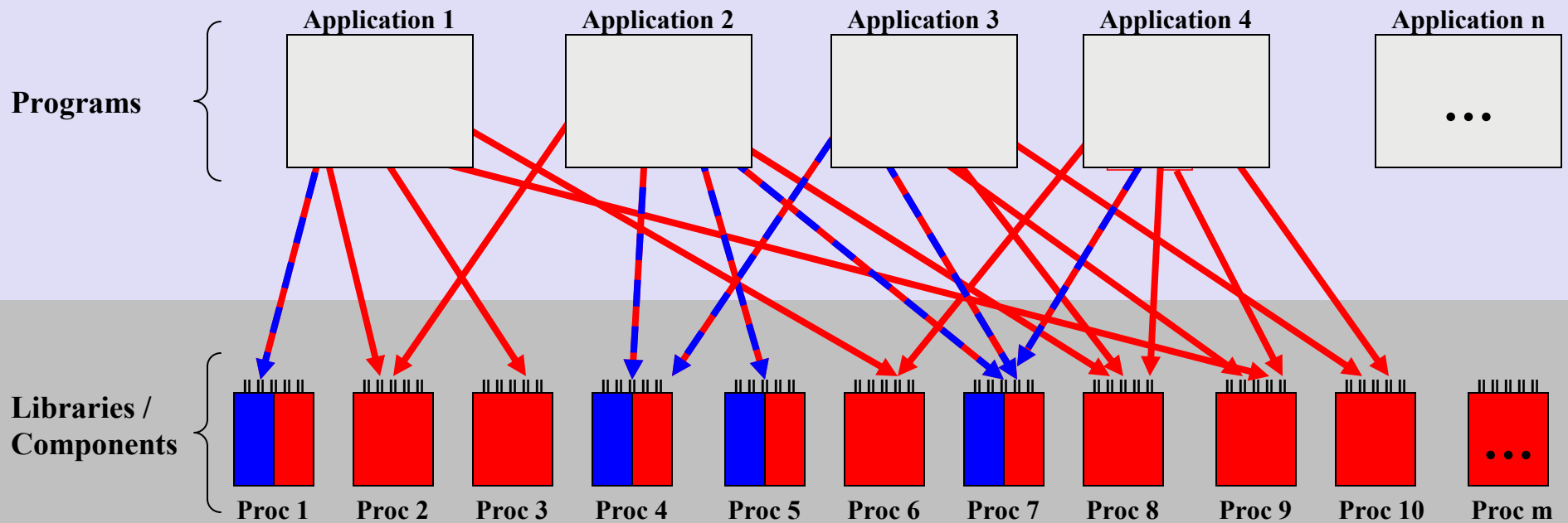
Generic marshalling

`xdrmem_putlong`

Buffer write + overflow check

`htonl`

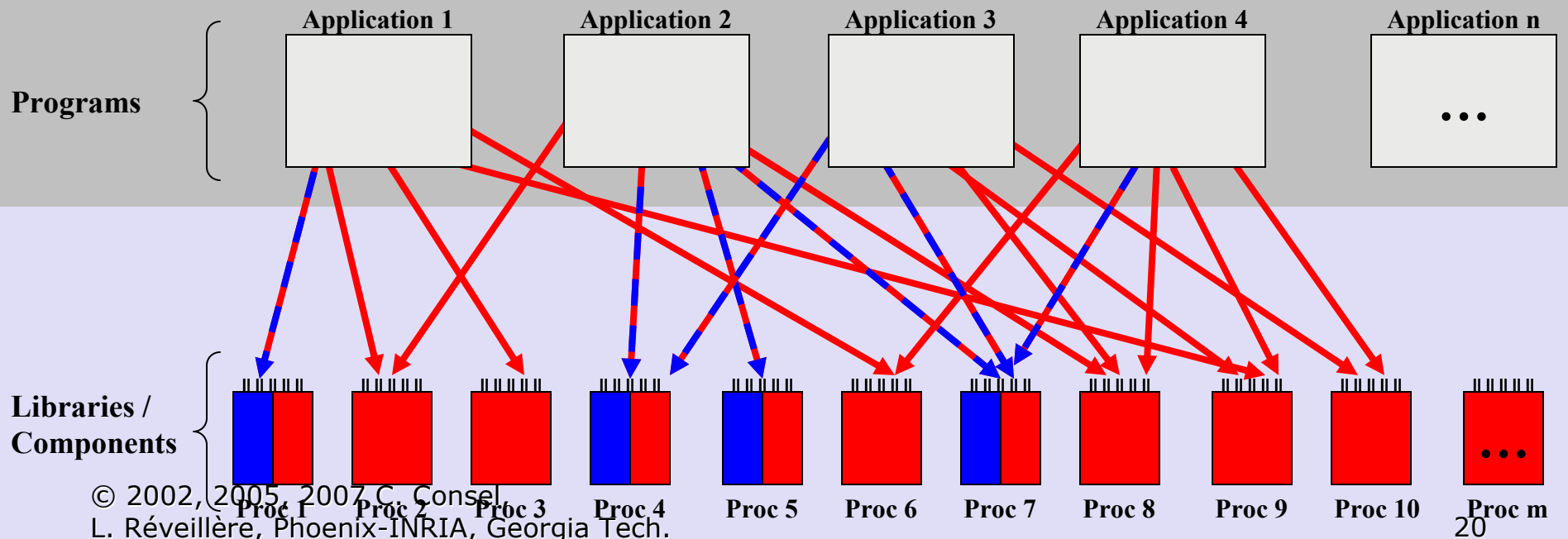
Byte ordering



# Program Family: Applications

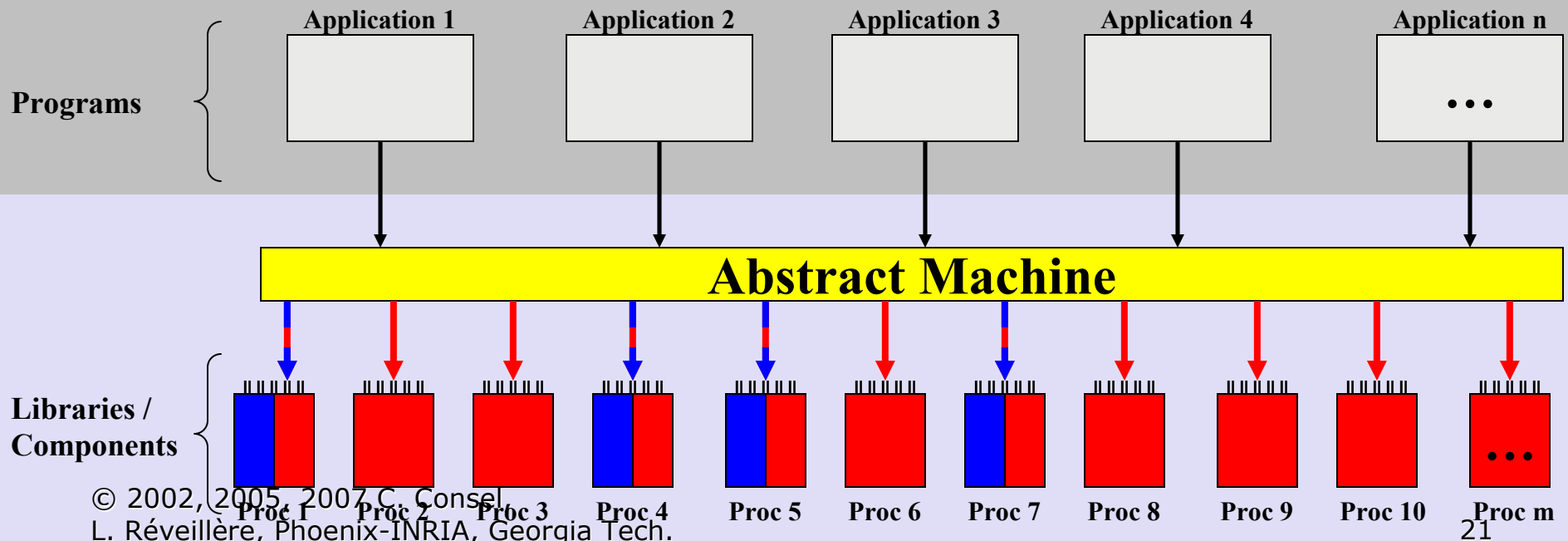
- ◆ General-purpose language
- ◆ Lack of software architecture

- Prologues/epilogues
  - Prepare invocations
  - Maintain state
  - Perform checks
- Results: complex, repetitive, expertise required



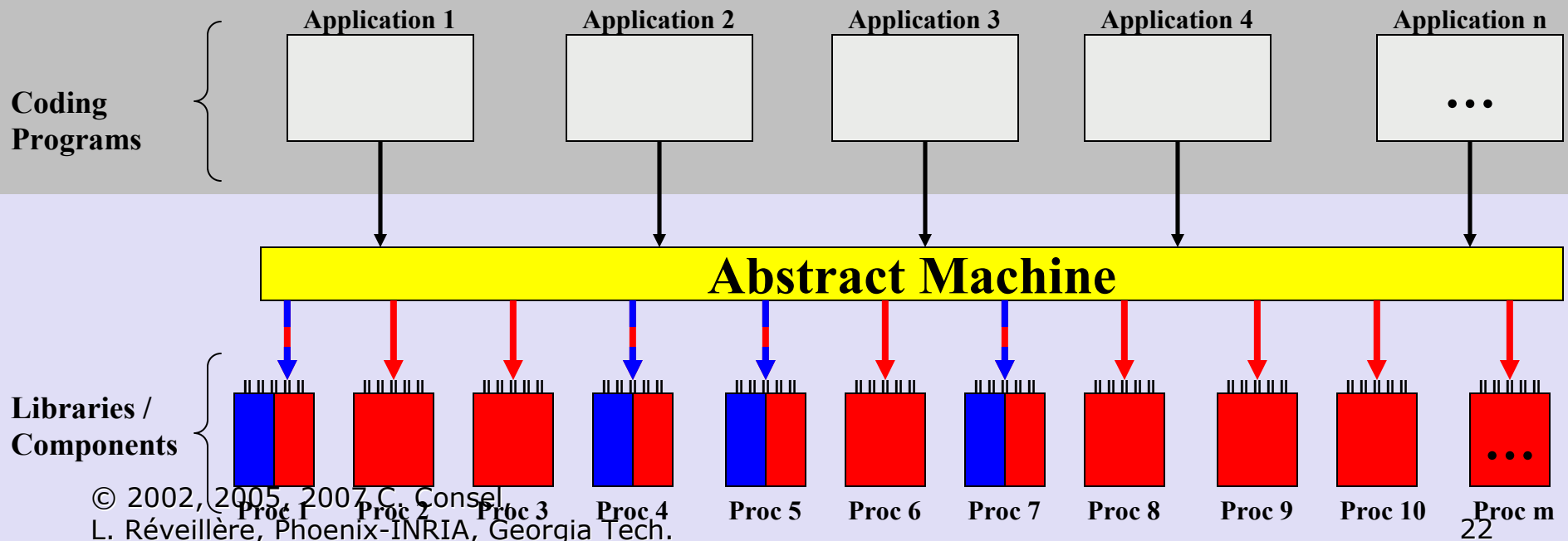
# Applications: *Need for Abstract Machine* (intermediate step)

- ◆ Interface to libraries
- ◆ Explicit run-time model
  - State definition and operations
  - Dynamic checks
  - Operation encapsulation
  - Operation combination
- More concise
- More usable
- More explicit expertise
- Abstraction layer



# Coding Applications: *Need for Abstract Machine* (intermediate step)

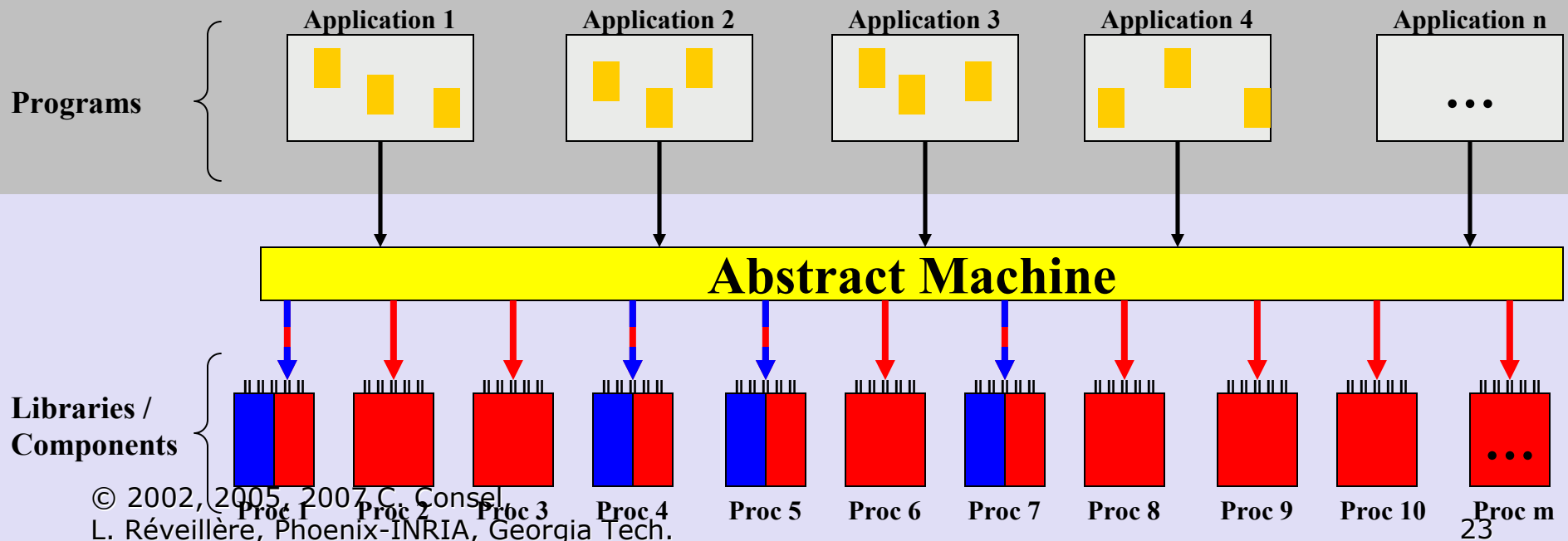
- More concise
  - More usable
  - More explicit expertise
  - Abstraction layer
- ◆ **XDR Data Coding Machine**
    - Instruction set: XDR library
    - State: XDR structure
    - Parameterized instruction set



# Program Family: Applications

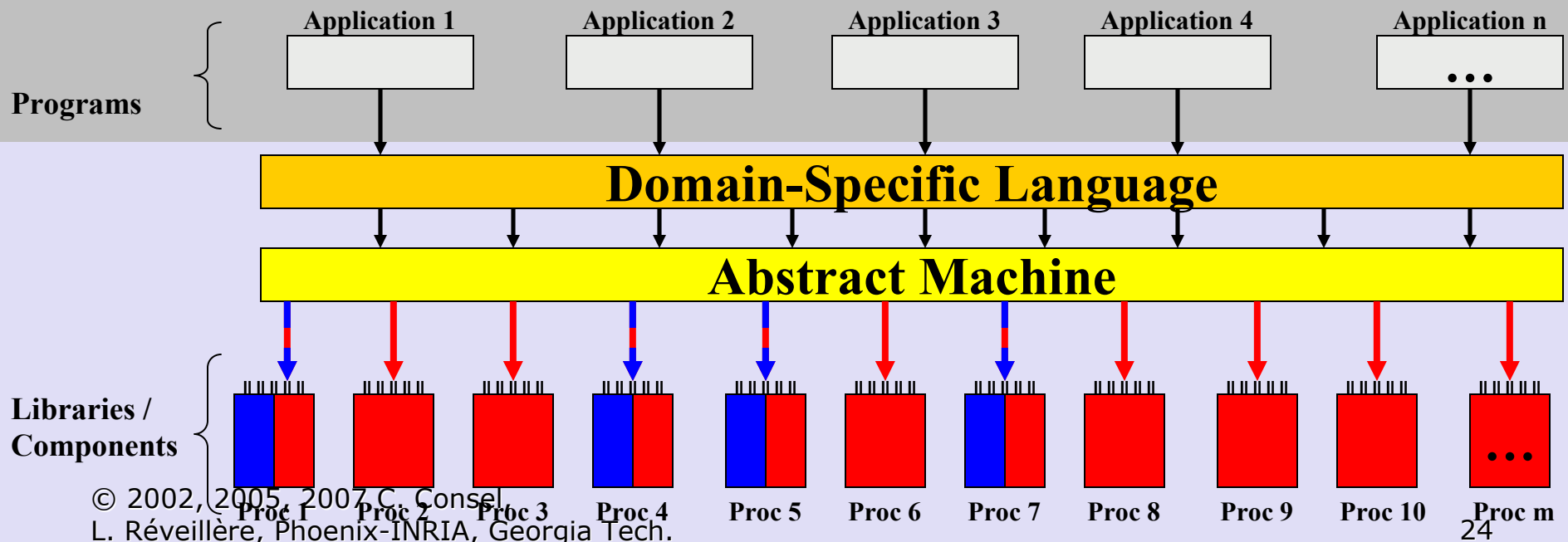
- ◆ General-purpose language

- For each call
  - Prologues/epilogues (still)
- For each call sequence
  - Repetitive/*required* call patterns
  - Repetitive/*required* checks



# Applications: *Need for Domain-Specific Language*

- ◆ Interface abstract machine
- ◆ Designate a family member
- ◆ Domain specific constructs (abstractions/notations)
- ◆ Restricted/enriched semantics
- Re-use
- Expertise
- Conciseness
- Verification
- Productivity
- Performance





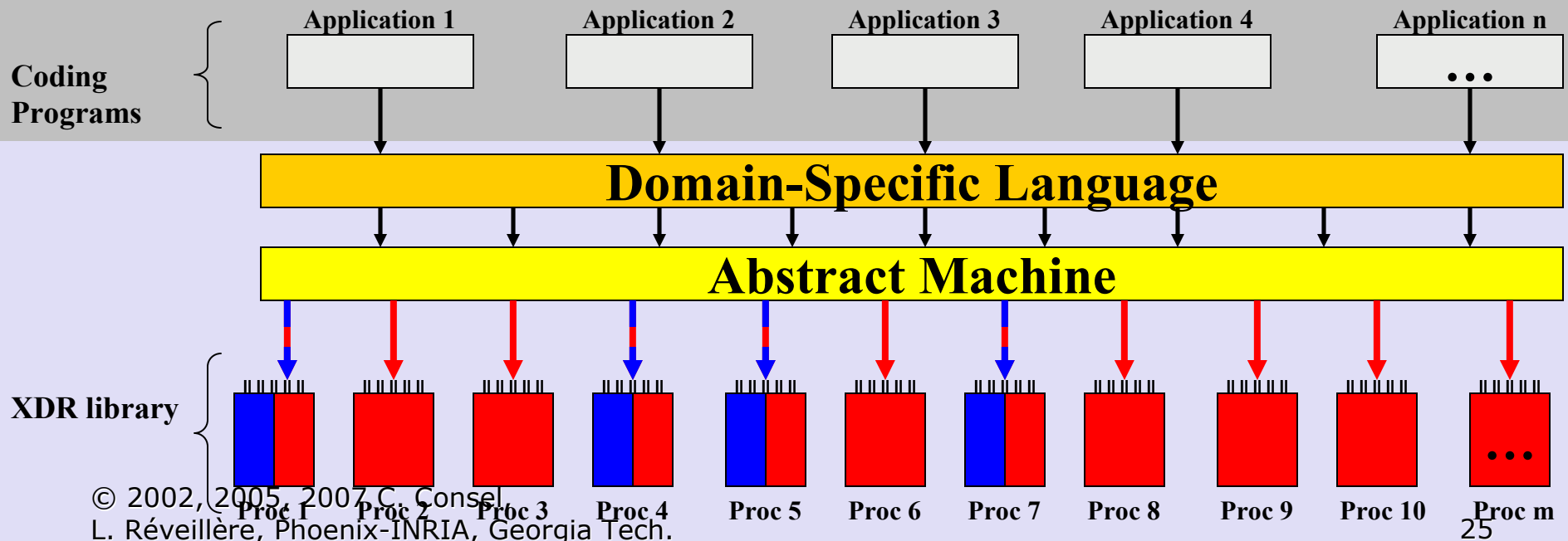
# Key concept:

## *Data Structure Descriptions*

### *XDR language*

```
union result switch (int error) {  
  case 0:  
    char data[MAX_SIZE];  
  default:  
    void;  
}
```

- Re-use
- Expertise
- Conciseness
- Verification
- Productivity
- Performance



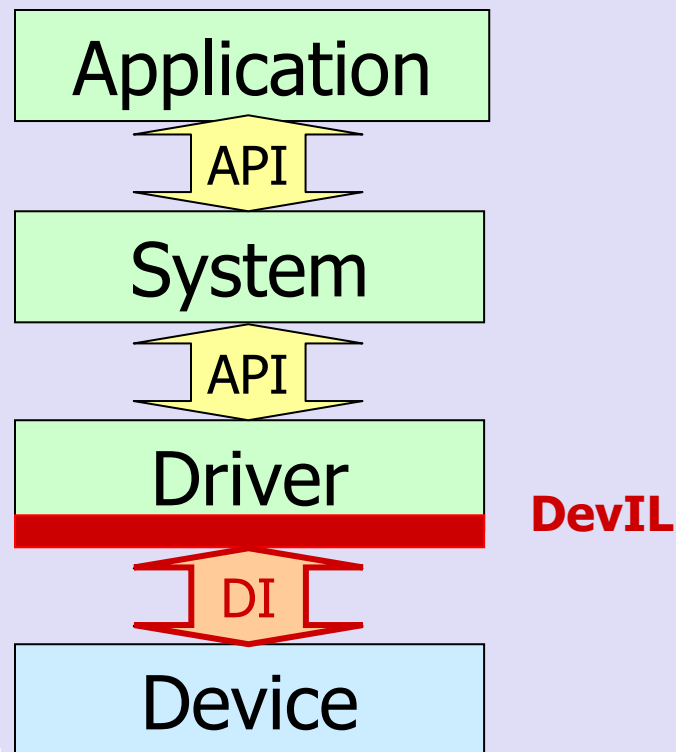
# Assessing a DSL

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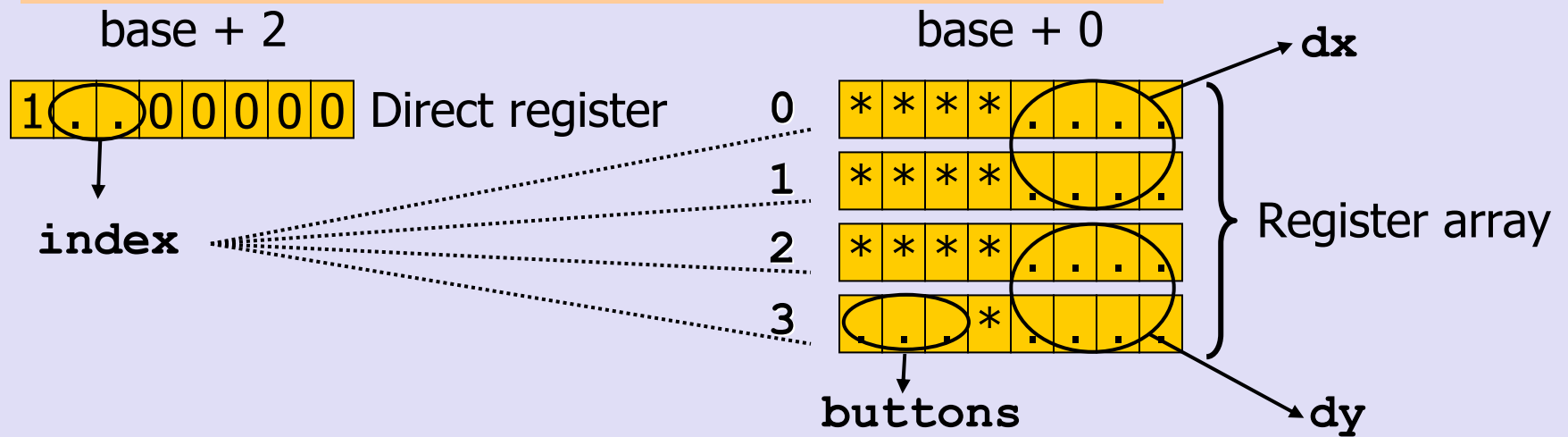
- ◆ Existing program family
  - (Almost) objective comparisons:
    - ◆ Conciseness
    - ◆ Robustness
    - ◆ Performance
  - (Likely) subjective comparisons:  
Readability, maintainability, usability

# A Program Family with our DSL: Device Driver Interface

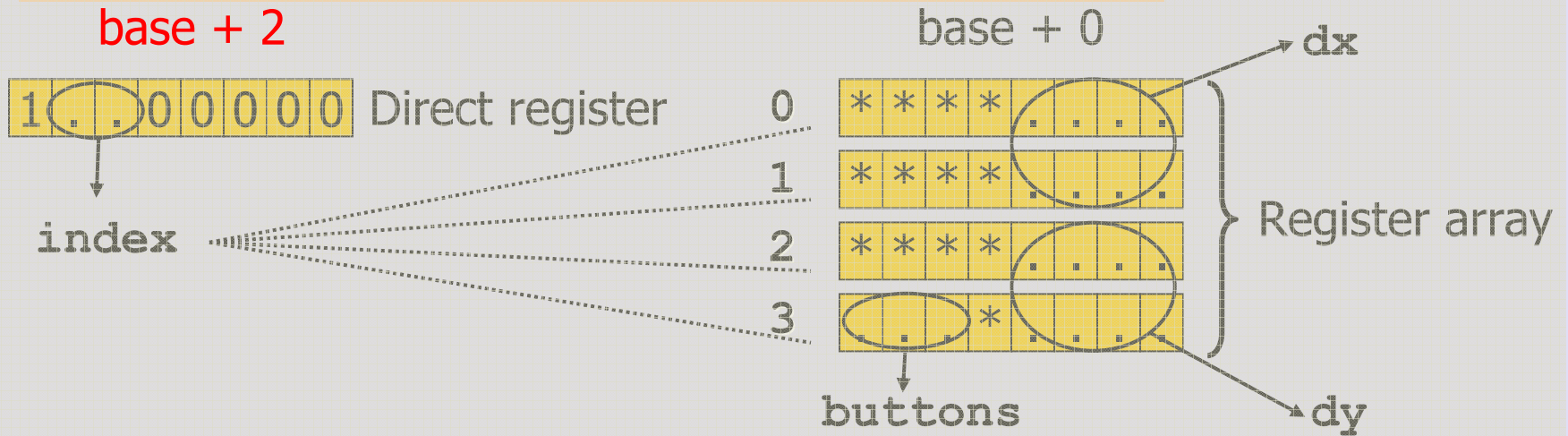
- ◆ Commonalities: API, bit operations...
- ◆ Variations: parameters, registers...



# Driver Example : Linux mouse



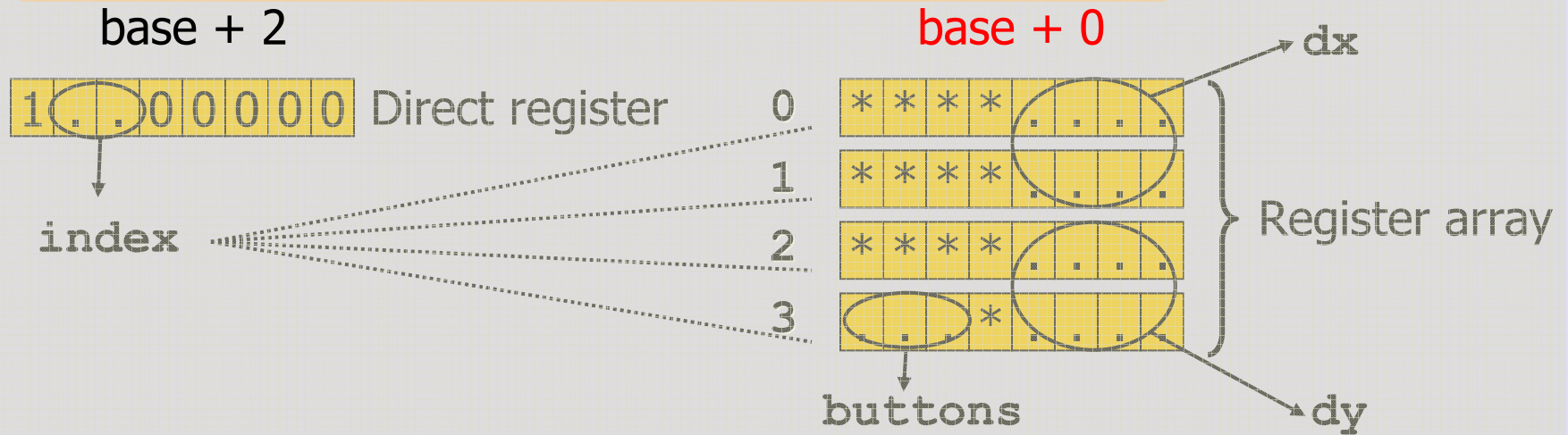
# Driver Example : Linux mouse



(busmouse.h)

```
#define MSE_CONTROL_PORT    0x23e
```

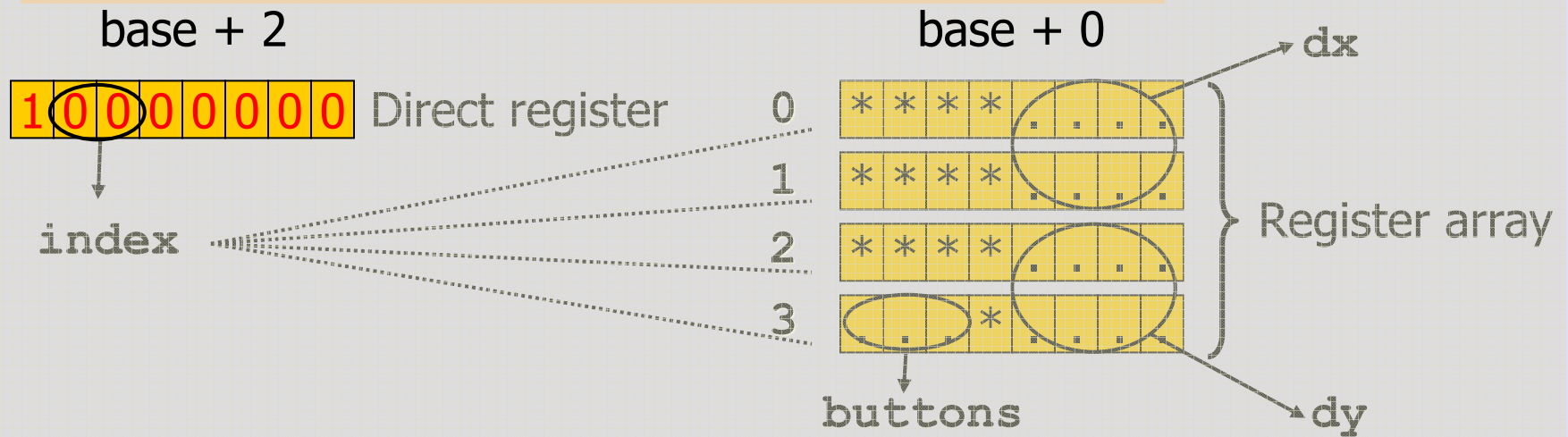
# Driver Example : Linux mouse



(busmouse.h)

```
#define MSE_CONTROL_PORT      0x23e
#define MSE_DATA_PORT        0x23c
```

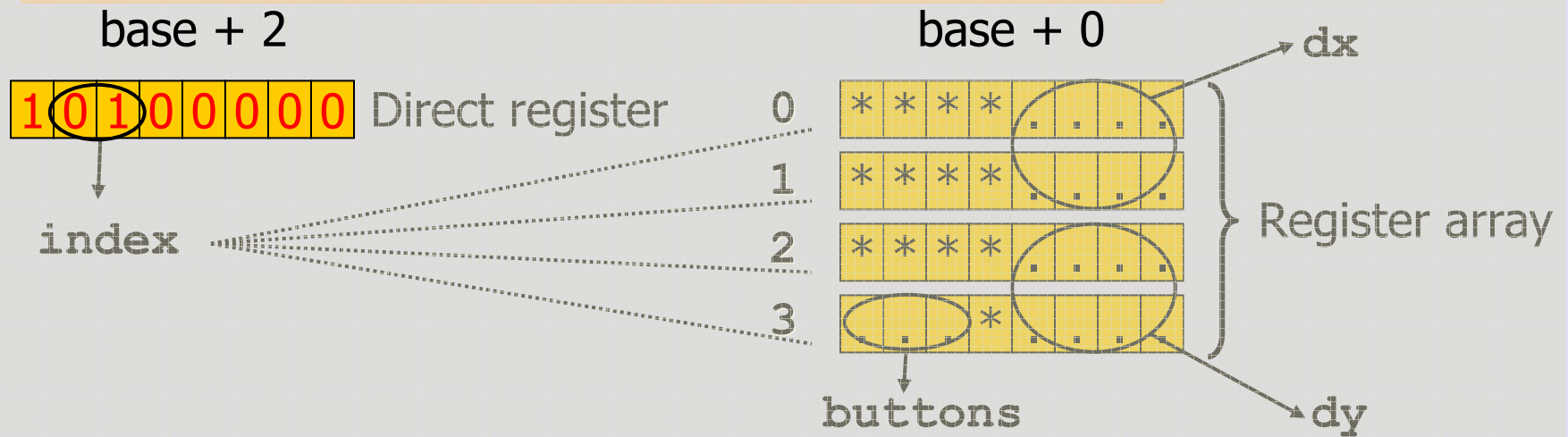
# Driver Example : Linux mouse



(busmouse.h)

```
#define MSE_DATA_PORT          0x23e
#define MSE_CONTROL_PORT      0x23c
#define MSE_READ_X_LOW        0x80
```

# Driver Example : Linux mouse

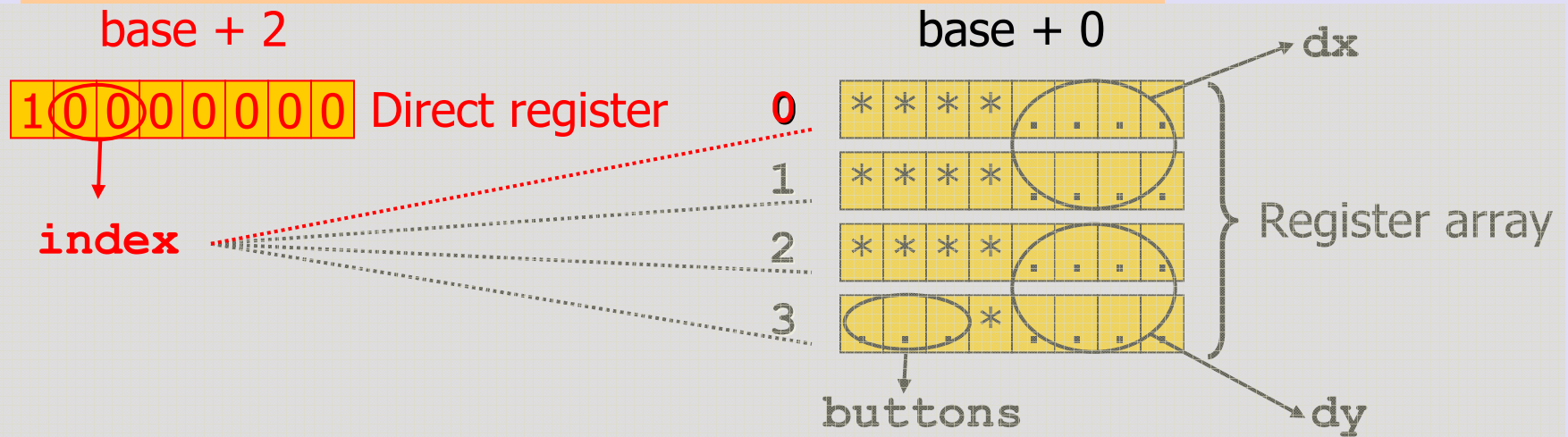


(busmouse.h)

```
#define MSE_DATA_PORT          0x23e
#define MSE_CONTROL_PORT      0x23c
#define MSE_READ_X_LOW        0x80
#define MSE_READ_X_HIGH       0xa0
```



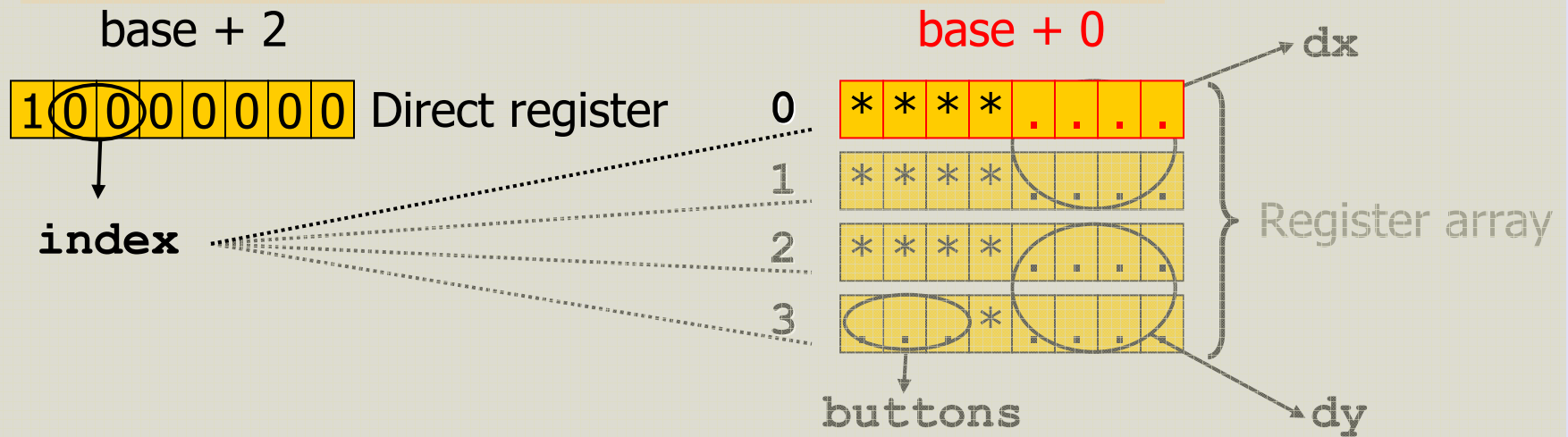
# Driver Example : Linux mouse



(busmouse.c)

```
outb(MSE_READ_X_LOW, MSE_CONTROL_PORT);
```

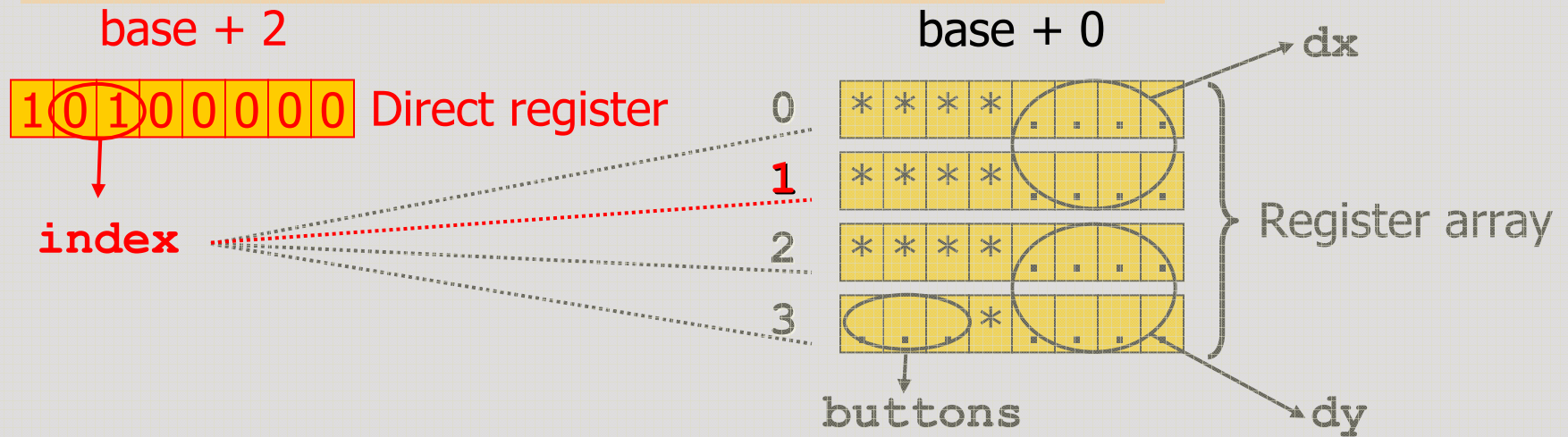
# Driver Example : Linux mouse



(busmouse.c)

```
outb(MSE_READ_X_LOW, MSE_CONTROL_PORT);  
dx = (inb(MSE_DATA_PORT) & 0xf);
```

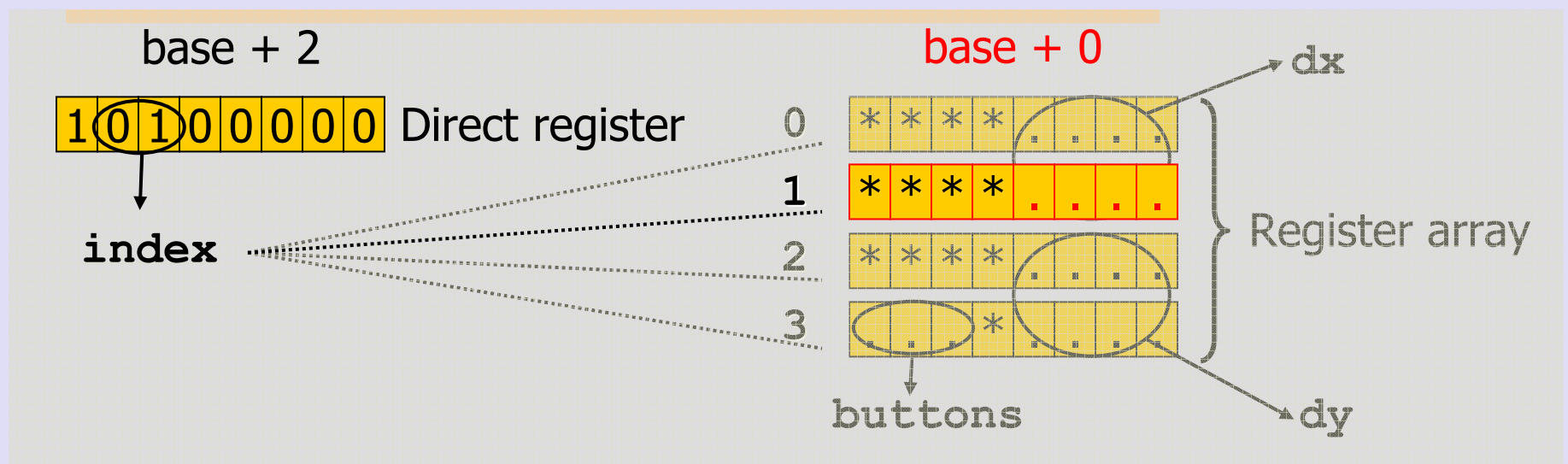
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```

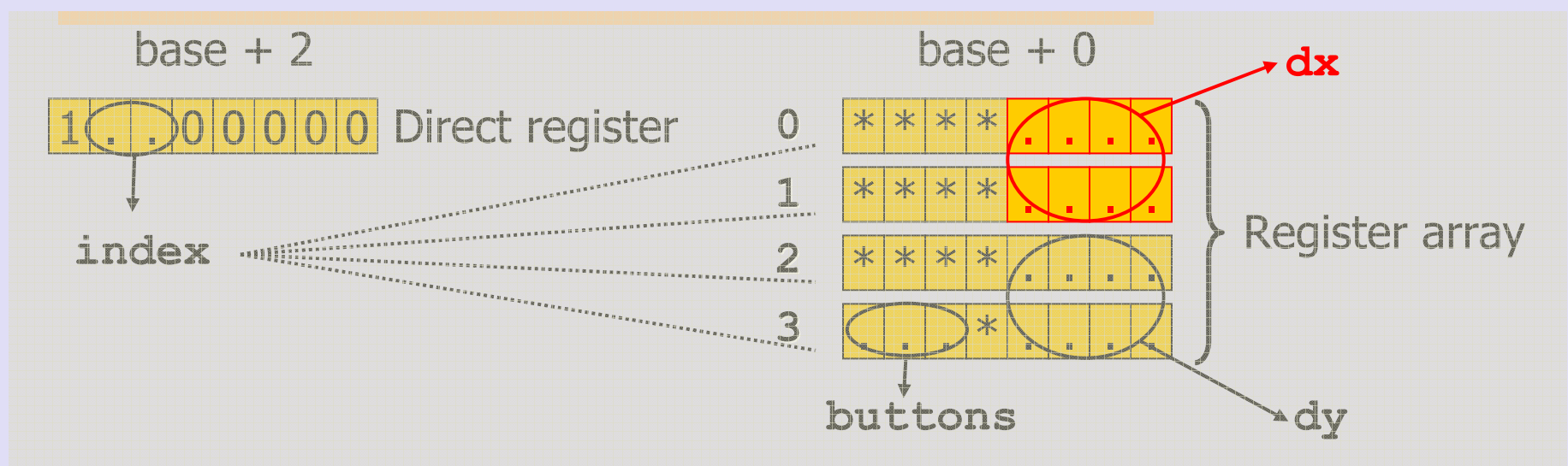
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dx = (inb(MSE_DATA_PORT) & 0xf);  
outb(MSE_READ_X_HIGH, MSE_CONTROL_PORT);  
dx |= (inb(MSE_DATA_PORT) & 0xf) << 4;
```

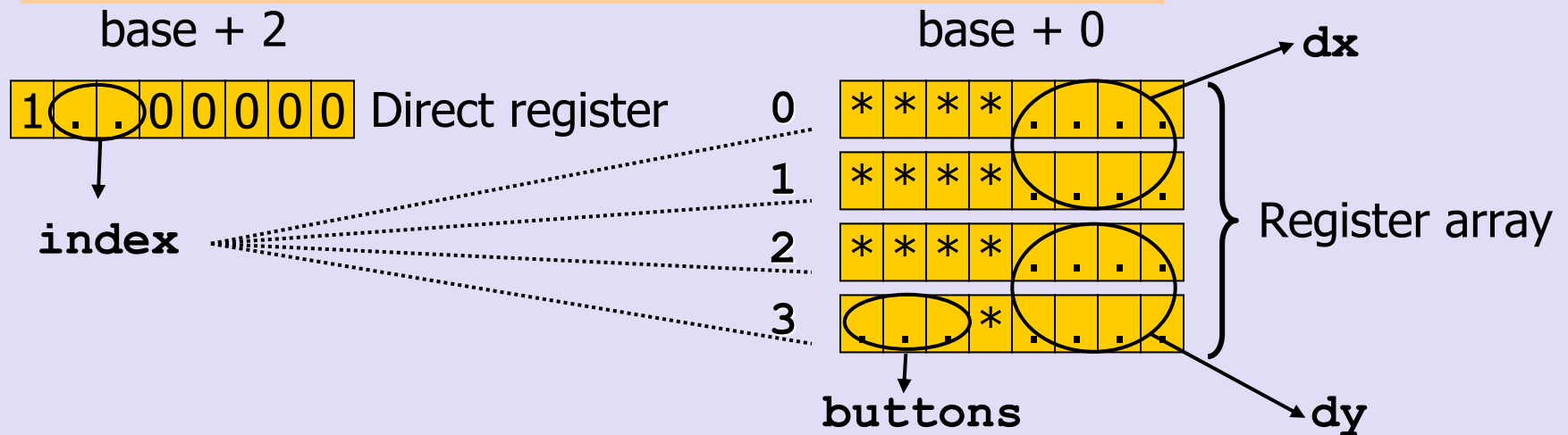
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```

# Driver Example : Linux mouse

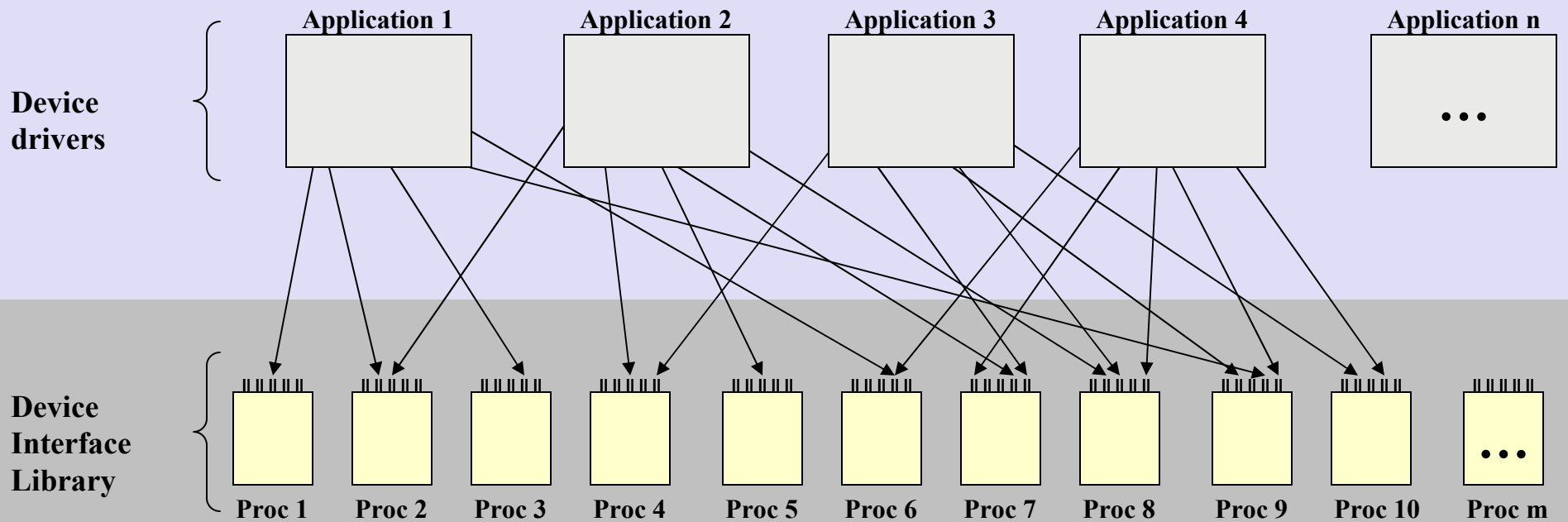


(busmouse.c)

```
outb(MSE_READ_X_LOW, MSE_CONTROL_PORT);  
dx = (inb(MSE_DATA_PORT) & 0xf);  
outb(MSE_READ_X_HIGH, MSE_CONTROL_PORT);  
dx |= (inb(MSE_DATA_PORT) & 0xf) << 4;  
...
```

# Family of Layers: Device Interface Library

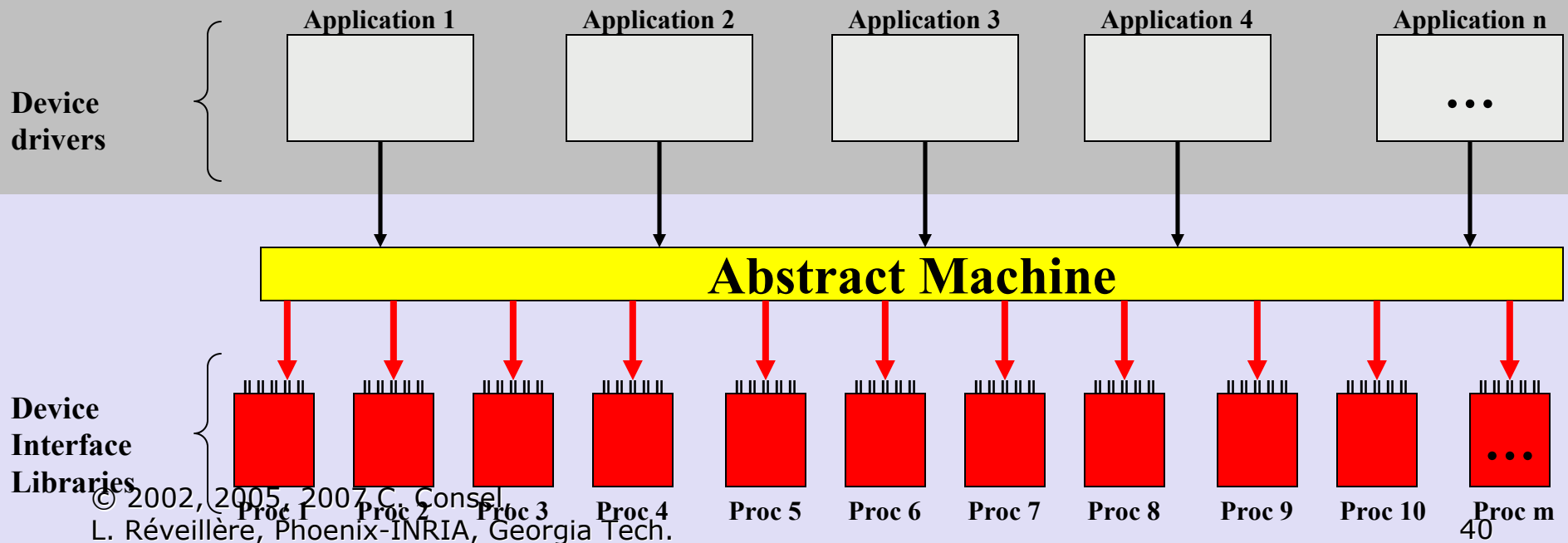
- ◆ Bit operations
- ◆ Input/output operations
- ◆ Low level
- ◆ Fine grained
- ◆ Directly mapped into machine instructions
- ◆ Implicit state



# Applications:

## *Device Interface Abstract Machine*

- ◆ Bit extraction operations
- ◆ Bit concatenation operation
- ◆ Input/output abstract operations
- ◆ Input/output checks
- ◆ Dedicated set of bit instructions
- ◆ Higher level
- ◆ More concise
- ◆ More robust

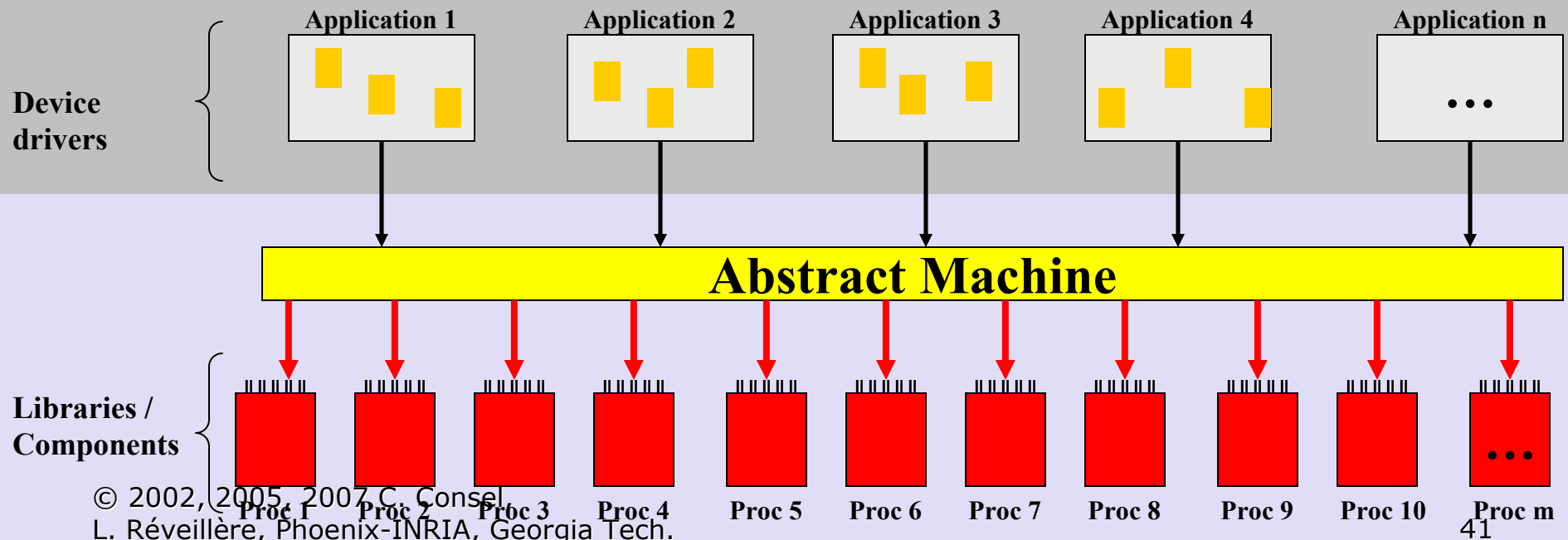




# Program Family: Applications

## ◆ General-purpose language

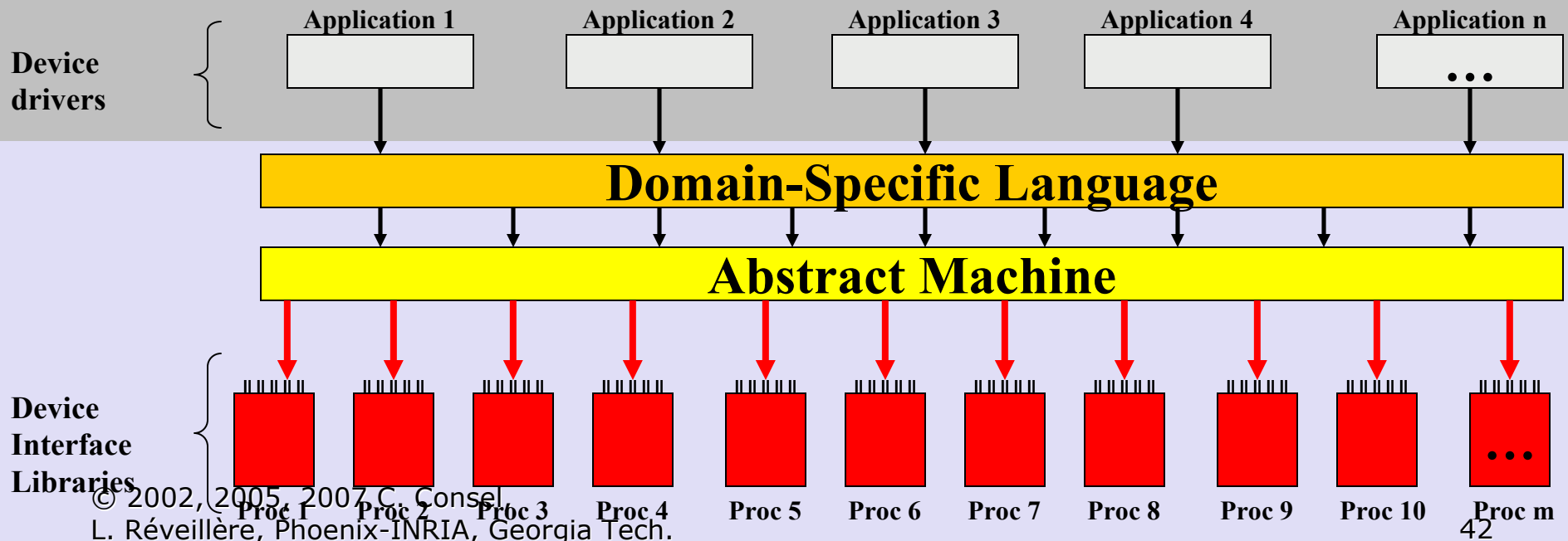
- General purpose glue for device communications
- Checks subject to programmer's rigor
- Repetitive communication sequences



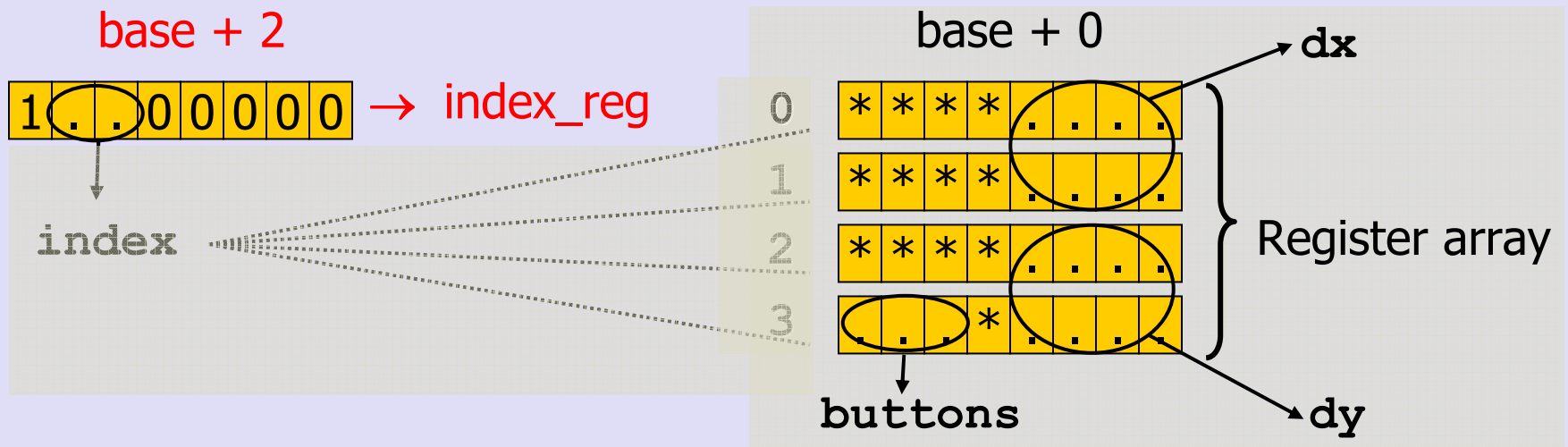
# Key concept:

## *Layered Device Interface*

- ◆ Port
  - Communication point
  - Different media: I/O, memory mapped
- ◆ Register
  - Repository of data in the device
  - indexed / paged, read / write, size, mask, ...
- ◆ Variable
  - Collection of register fragments
  - Semantic value

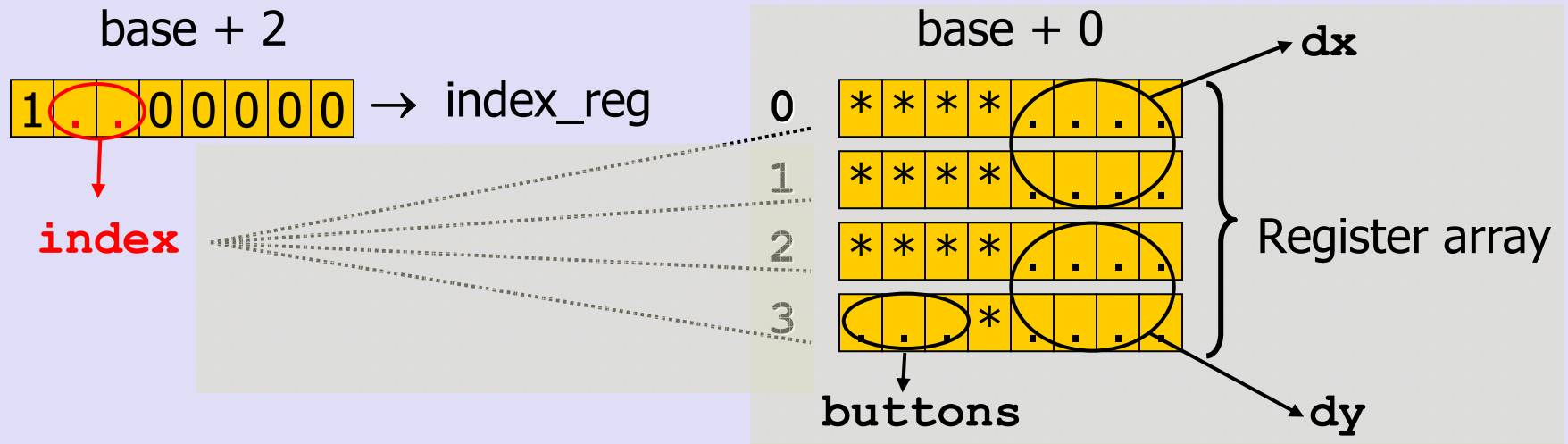


# Direct Access to Register



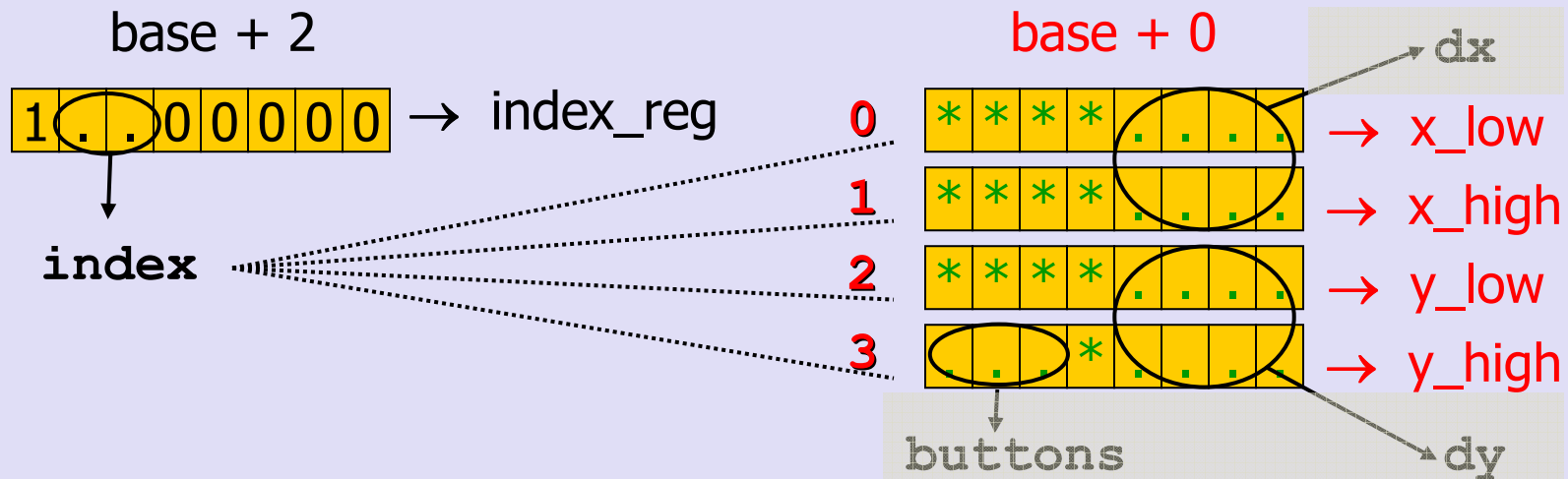
```
...  
register index_reg = write base@2, mask '1..00000': bit[8];
```

# Private Variable



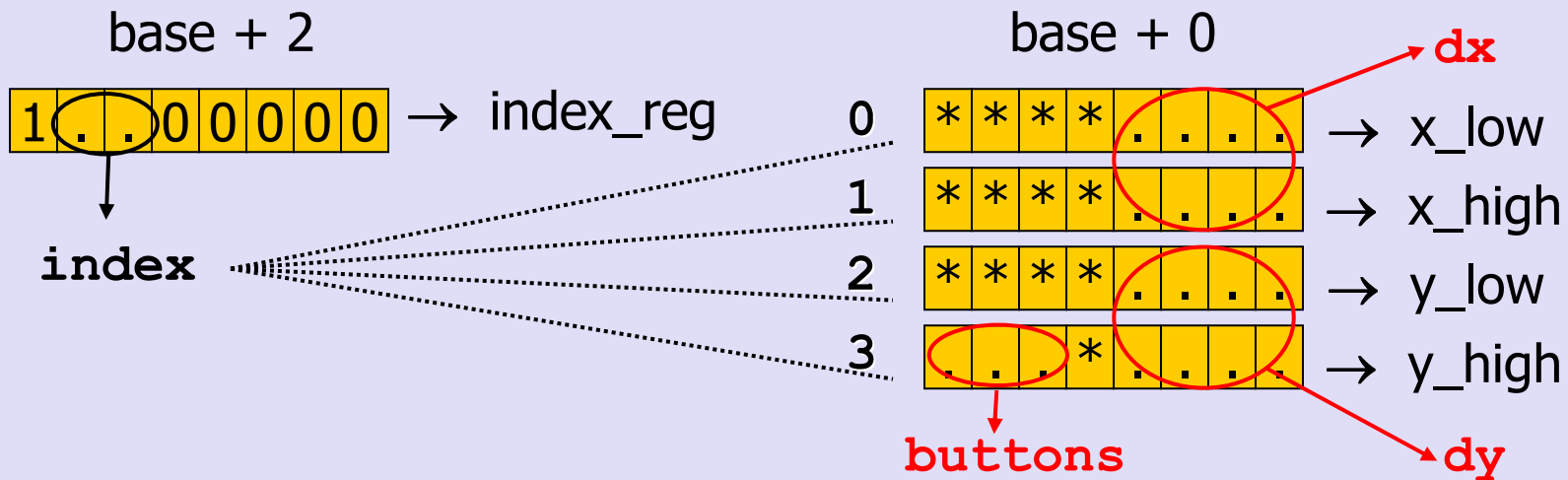
```
...  
register index_reg = write base@2, mask '1..00000' : bit[8];  
private variable index = index_reg[6..5] : int(2);
```

# Indexed Registers



```
...
register index_reg = write base@2, mask '1..00000' : bit[8];
private variable index = index_reg[6..5] : int(2);
...
register x_low  = read base@0, pre {index = 0}, mask '****....' : bit[8];
register x_high = read base@0, pre {index = 1}, mask '****....' : bit[8];
...
```

# Variables



```
...
register index_reg = write base@2, mask '1..00000' : bit[8];
private variable index = index_reg[6..5] : int(2);
...
register x_low = read base@0, pre {index = 0}, mask '****....' : bit[8];
register x_high = read base@0, pre {index = 1}, mask '****....' : bit[8];
...
variable dx = x_high[3..0] # x_low[3..0], volatile : signed int(8);
variable dy = y_high[3..0] # y_low[3..0], volatile : signed int(8);
```

# Comparison

```
outb(MSE_READ_X_LOW, MSE_CONTROL_PORT);
dx = (inb(MSE_DATA_PORT) & 0xf);
outb(MSE_READ_X_HIGH, MSE_CONTROL_PORT);
dx |= (inb(MSE_DATA_PORT) & 0xf) << 4;
outb(MSE_READ_Y_LOW, MSE_CONTROL_PORT);
dy = (inb(MSE_DATA_PORT) & 0xf);
outb(MSE_READ_Y_HIGH, MSE_CONTROL_PORT);
buttons = inb(MSE_DATA_PORT);
dy |= (buttons & 0xf) << 4;
```

Existing  
code

Devil

```
variable dx = x_high[3..0] # x_low[3..0], volatile
            : signed int(8);
variable dy = y_high[3..0] # y_low[3..0], volatile
            : signed int(8);
```

# Talk Outline

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- ◆ Introduction to Domain Specific Languages (DSL)
- ◆ **Overview of SIP**
- ◆ SPL: A DSL for communication services
- ◆ Properties of SPL
- ◆ Summary

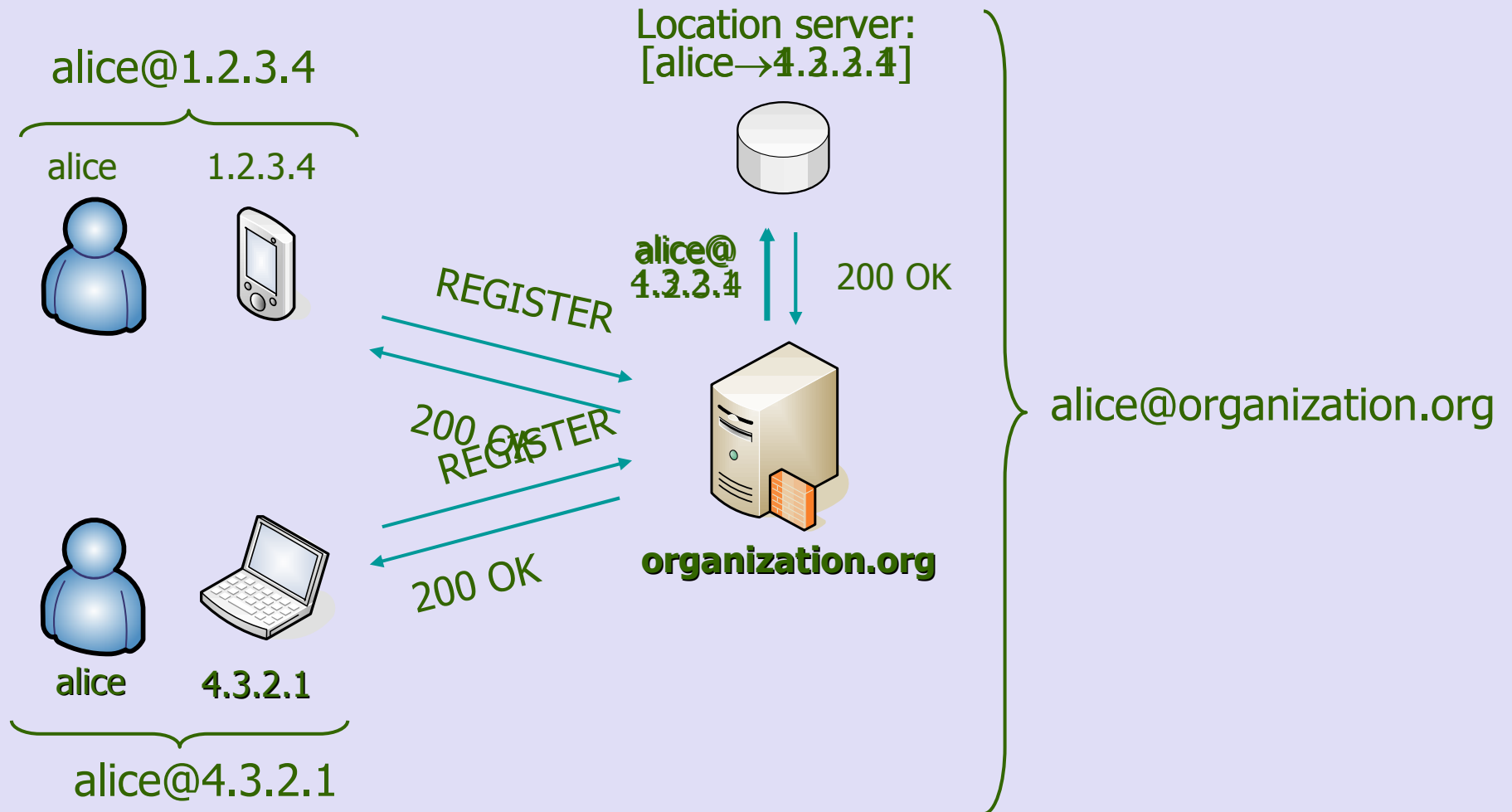


# What is SIP?

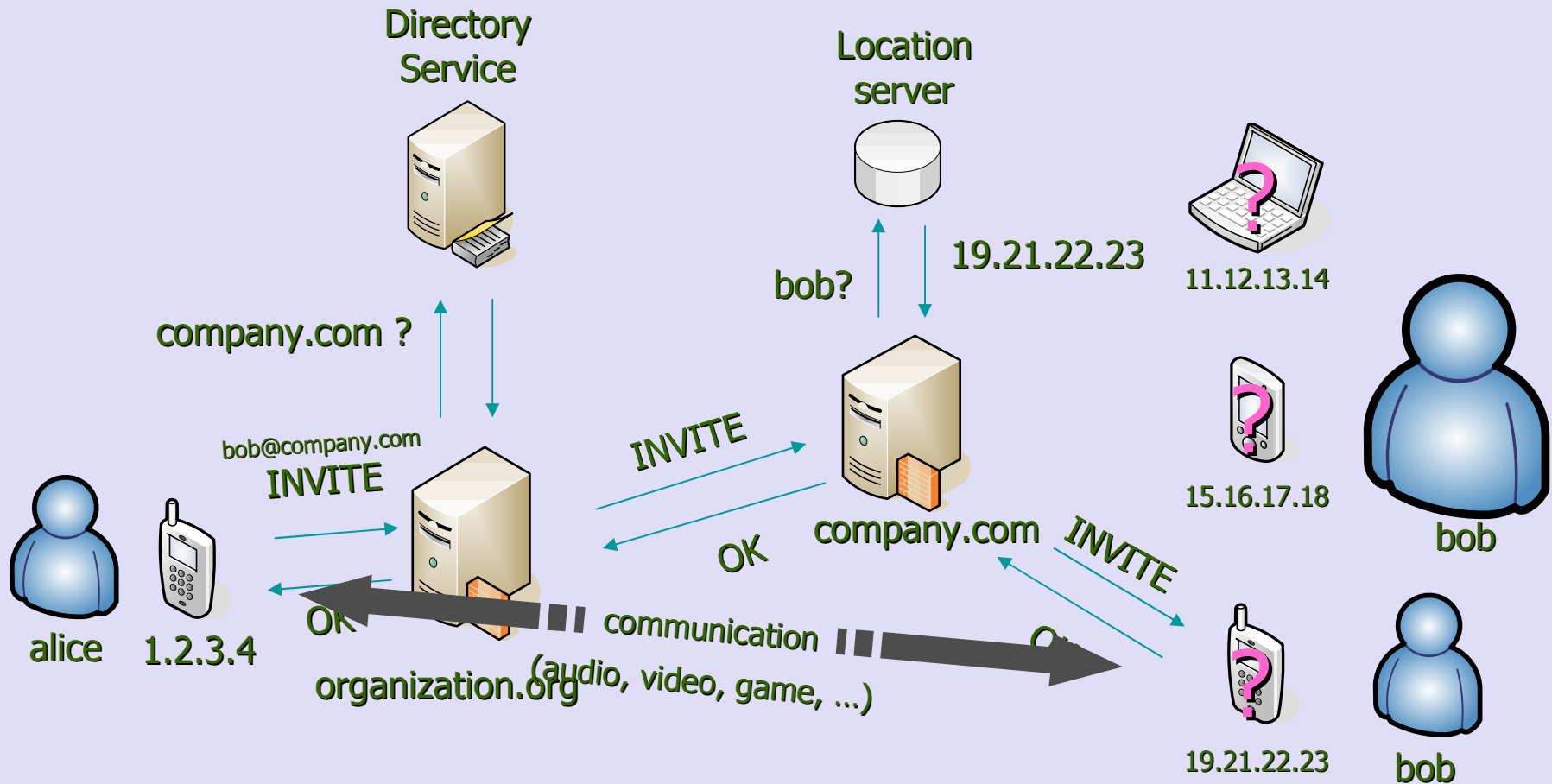
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- ◆ SIP: Session Initiation Protocol
  - Voice over IP and 3G mobile phones
  - Standardized by IETF, adopted by ITU
- ◆ Protocol for creating, modifying and terminating sessions between devices
- ◆ Based on a client/server model
- ◆ Mobility: SIP URI

# SIP: Mobility

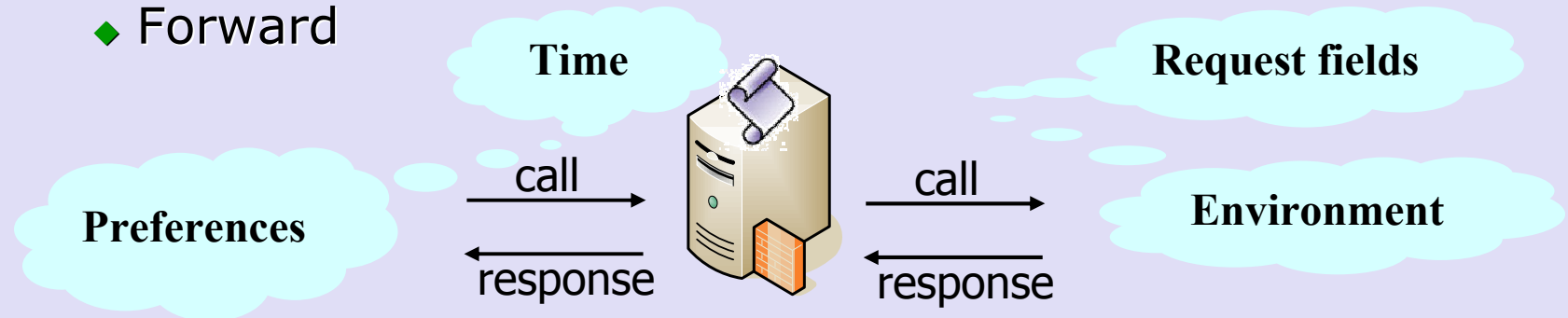


# SIP: Typical Call Flow



# Programming Telephony Services

- ◆ A service processes SIP messages (incoming/outgoing) on behalf of a person or a group of persons
  - Service logic
  - Signalling actions
    - ◆ Reject
    - ◆ Forward

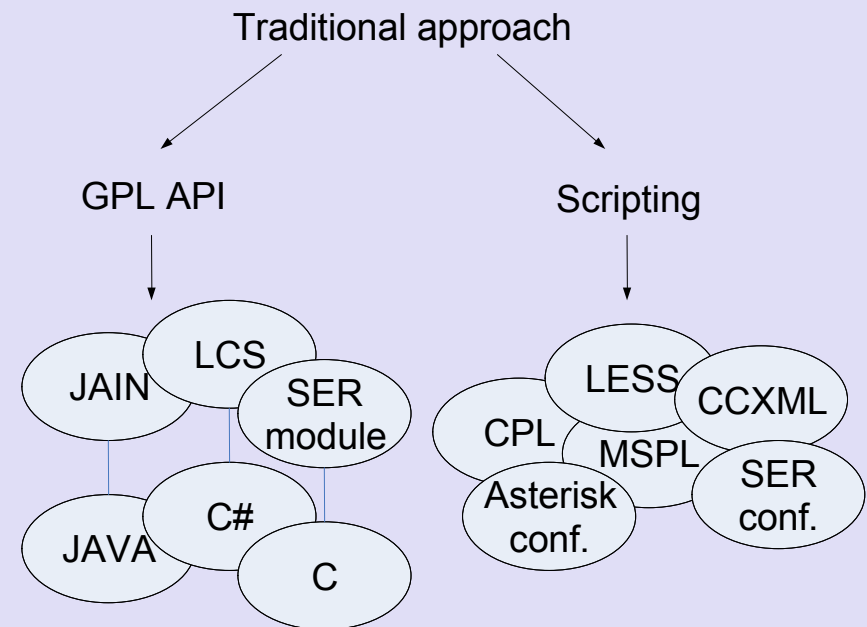


# Existing Solutions

## ◆ Use General Purpose Languages:

- ◆ Microsoft SIP APIs – C#
- ◆ SIP Express Router – C-like and C
- ◆ JAIN SIP – Java

## ◆ Requires extensive knowledge of existing protocols



Expressive, but unsafe

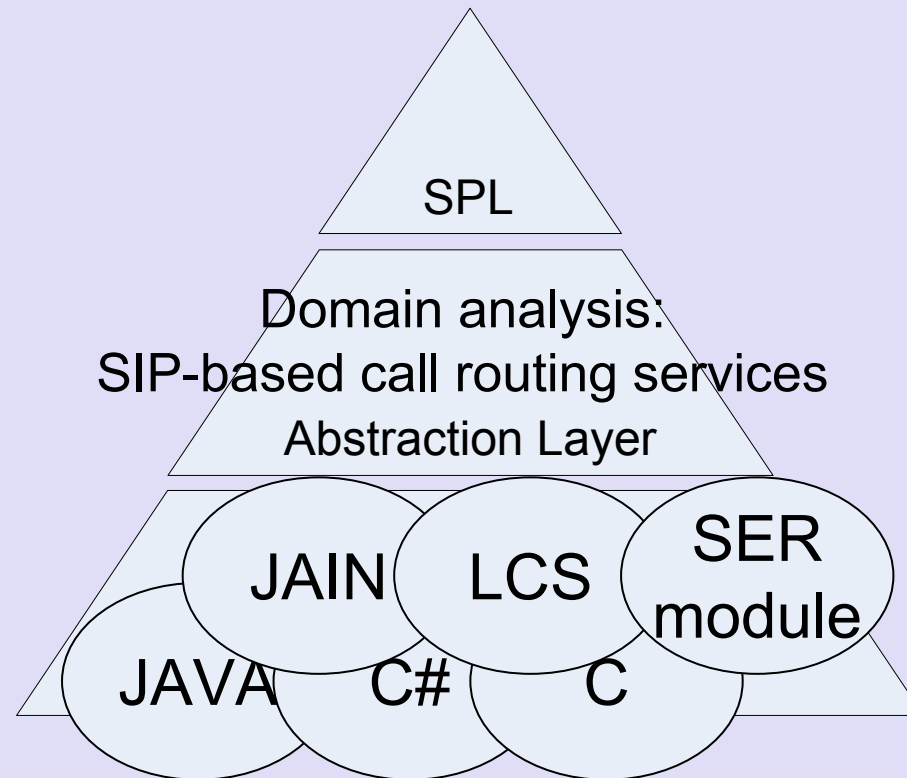
# Why a DSL for Telephony Services?

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- ◆ Telephony services are a rich domain
- ◆ SIP is a rich protocol
- ◆ SIP has many companion protocols
- ◆ Programming services requires expertise in multiple areas
  - Protocols
  - Networking
  - Platform
- ◆ Few trained programmers
- ◆ Reliability is critical
- ◆ Performance is critical

# DSL Based Approach

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# Example

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- ◆ Abstract service specification
  - When registering
    - ◆ Set a counter to 0
  - When a call is received
    - ◆ If the callee rejects the call :
      - Forward the call to the secretary
    - ◆ Otherwise
      - Increment the counter
  - When un-registering
    - ◆ Log the value of the counter



# Example: JAIN SIP Code

```
public class Counter implements SipListener {

    public void processRequest(RequestEvent requestEvent) {
        try {
            Request rq_request = requestEvent.getRequest();
            SipProvider rq_sipProvider = (SipProvider) requestEvent.getSource();
            String method = rq_request.getMethod();
            if (method.equals(Request.REGISTER)) {
                if (!registrar.hasExpiresZero(rq_request)) {
                    if (!registrar.hasRegistration(rq_request)) {
                        State state = new State();
                        int ident = env.getId(rq_request);
                        env.setEnv(ident, state);
                        state.count = 0;
                        rq_sipProvider.sendRequest(rq_request);
                        return;
                    } else {
                        rq_sipProvider.sendRequest(rq_request);
                        return;
                    }
                } else if (lib.registrar.hasRegistration(rq_request)) {
                    int ident = env.getId(rq_request);
                    State state = (State)env.getEnv(ident);
                    log(state.count);
                    env.delEnv(ident);
                    rq_sipProvider.sendRequest(rq_request);
                    return;
                } else {
                    rq_sipProvider.sendRequest(rq_request);
                    return;
                }
            } else if (method.equals(Request.INVITE)) {
                ClientTransaction ct = rq_sipProvider.getNewClientTransaction(rq_request);
                ct.sendRequest();
                return;
            } else {
                rq_sipProvider.sendRequest(rq_request);
                return;
            }
        } catch (Exception ex) { [...] }
    }
}
```

[...]

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L. Réveillère, Phoenix-INRIA, Georgia Tech.

```
[...]

    public void processResponse(ResponseEvent responseEvent) {
        try {
            Response rs_response = responseEvent.getResponse();
            SipProvider rs_sipProvider = (SipProvider) responseEvent.getSource();
            rs_responseCode = rs_response.getStatusCode();
            ClientTransaction ct = responseEvent.getClientTransaction();
            if ( clientTransaction != null ) {
                Request rs_request = ct.getRequest();
                String method = rs_request.getMethod();
                int ident = env.getId(rs_request);
                State state = (State)env.getEnv(ident);
                if (method.equals(Request.INVITE)) {
                    if (rs_responseCode >= 300) {
                        AddressFactory addressFactory = lib.getAddressFactory();
                        SipURI sipURI =
                            addressFactory.createSipURI("secretary","company.com");
                        rs_request.setRequestURI(sipURI);
                        int ident = env.getId(rs_request);
                        env.delEnv(ident);
                        sipProvider.sendRequest(rs_request);
                        return;
                    } else {
                        state.count++;
                        ServerTransaction st = rs_request.getNewServerTransaction();
                        st.sendResponse(rs_response);
                        return;
                    }
                } else {
                    rs_sipProvider.sendResponse(rs_response);
                    return;
                }
            } else {
                rs_sipProvider.sendResponse(rs_response);
                return;
            }
        } catch (Exception ex) { [...] }
    }
}
```

[...]

# Example: SPL Code

```
service Counter {
  processing {
    local void log (int);

    registration {
      int count;

      response outgoing REGISTER() {
        count = 0;
        return forward;
      }

      void unregister() {
        log (count);
      }

      dialog {
        response incoming INVITE() {
          response resp = forward;
          if (resp != /SUCCESS) {
            return forward 'sip:secretary@company.com';
          } else {
            count++;
            return resp;
          }
        }
      }
    }
  }
}
```

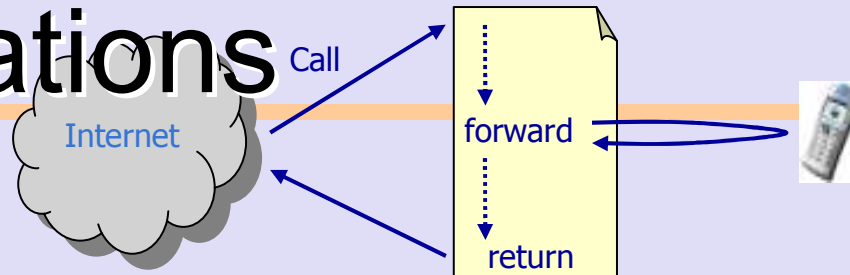
}}}}}  
© 2002, 2005, 2007 C. Consel,  
L. Réveillère, Phoenix-INRIA, Georgia Tech.

# SPL Components

---

- ◆ Event handlers and signaling operations
- ◆ Session
- ◆ Hierarchical sessions
  - Service
  - Registration
  - Dialog
- ◆ Inter-event control flow

# Event Handlers and Signaling Operations



## SPL Service

```
response incoming INVITE() {  
    [...]  
    response resp = forward;  
    if (resp == /ERROR) {  
        resp = forward 'sip:phoenix.secretary@inria.fr';  
    }  
    return resp;  
}
```

```
// Deny Service  
response incoming INVITE() {  
    return /ERROR/CLIENT/BUSY_HERE;  
}
```

# Session

---

- **INVITE**
- **BYE**
- **REINVITE**
- ...



Dialog = ID

```
uri caller;  
time start;  
...
```

# Session: The Dialog Session Example

```
dialog {
    uri caller;
    time start;

    response incoming INVITE() {
        caller = FROM;
        return forward;
    }

    void incoming ACK(){
        if(caller == 'sip:my.wife@home.fr')
            log("Personal call");
        start = getTime();
    }

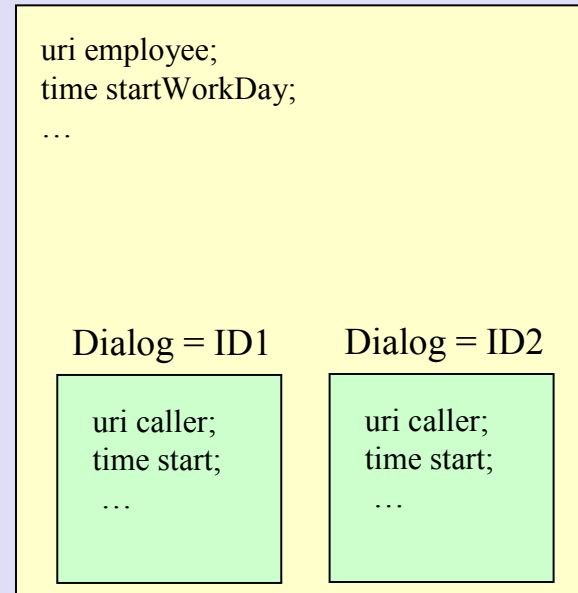
    response BYE() {
        string duration = time_to_string(getTime() - start);
        log("Call: "+ duration +" "+uri_to_string(caller));
        return forward;
    }
}
```

# Hierarchical Sessions

- **REGISTER**
- **REREGISTER**
- *unregister*

- **INVITE**
- **BYE**
- **REINVITE**
- ...

Registration = ID



# Hierarchical Sessions: Example

```
registration {
  uri employee;
  time startWorkDay;

  response outgoing REGISTER() {
    startWorkDay = getTime();
    employee = FROM;
    return forward;
  }

  void unregister() {
    string duration = time_to_string(getTime() - startWorkDay);
    log("WorkDay: "+ duration +" "+uri_to_string(employee));
    return;
  }

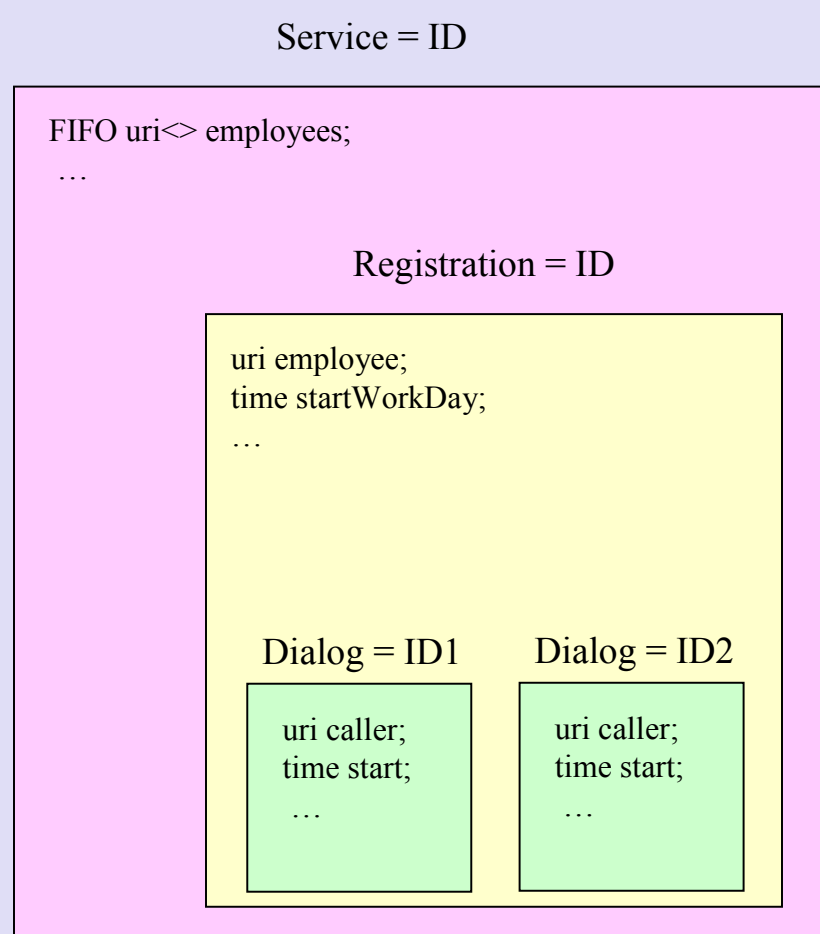
  dialog {
    uri caller; time start;

    ...
  }
}
```



# Hierarchical Sessions

- *deploy*
- *undeploy*
  - **REGISTER**
  - **REREGISTER**
  - *unregister*
- **INVITE**
  - **BYE**
  - **REINVITE**
  - ...



# Hierarchical Sessions: Example (cont'd)

```
service hotline {
  ...
  processing {
    uri<100> employees = <>;

    void deploy() {...}
    void undeploy() {...}

    registration {...

      response outgoing REGISTER() {
        startWorkDay = getTime();
        employee = FROM;
        push employees employee;
        return forward;
      }
      ...

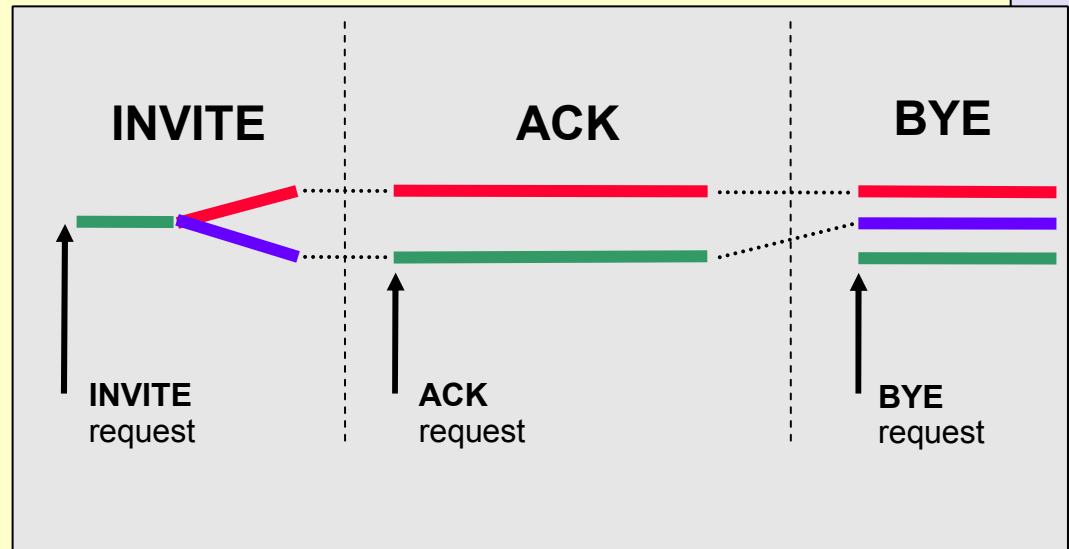
      dialog { ...
        response incoming INVITE() {
          return forward employees;
        }
      }
    }
  }
}
```

# Inter-Event Control Flow

```
dialog {
  response incoming INVITE() {
    response r;
    ...
    if (...) {
      ...
      return r branch hotline;
    }
    else {
      ...
      return r branch personal;
    }
  }

  void incoming ACK() {
    branch hotline {... }
    branch default {... }
  }

  response BYE() {
    branch hotline {... }
    branch personal {... }
    branch default {... }
  }
}
```



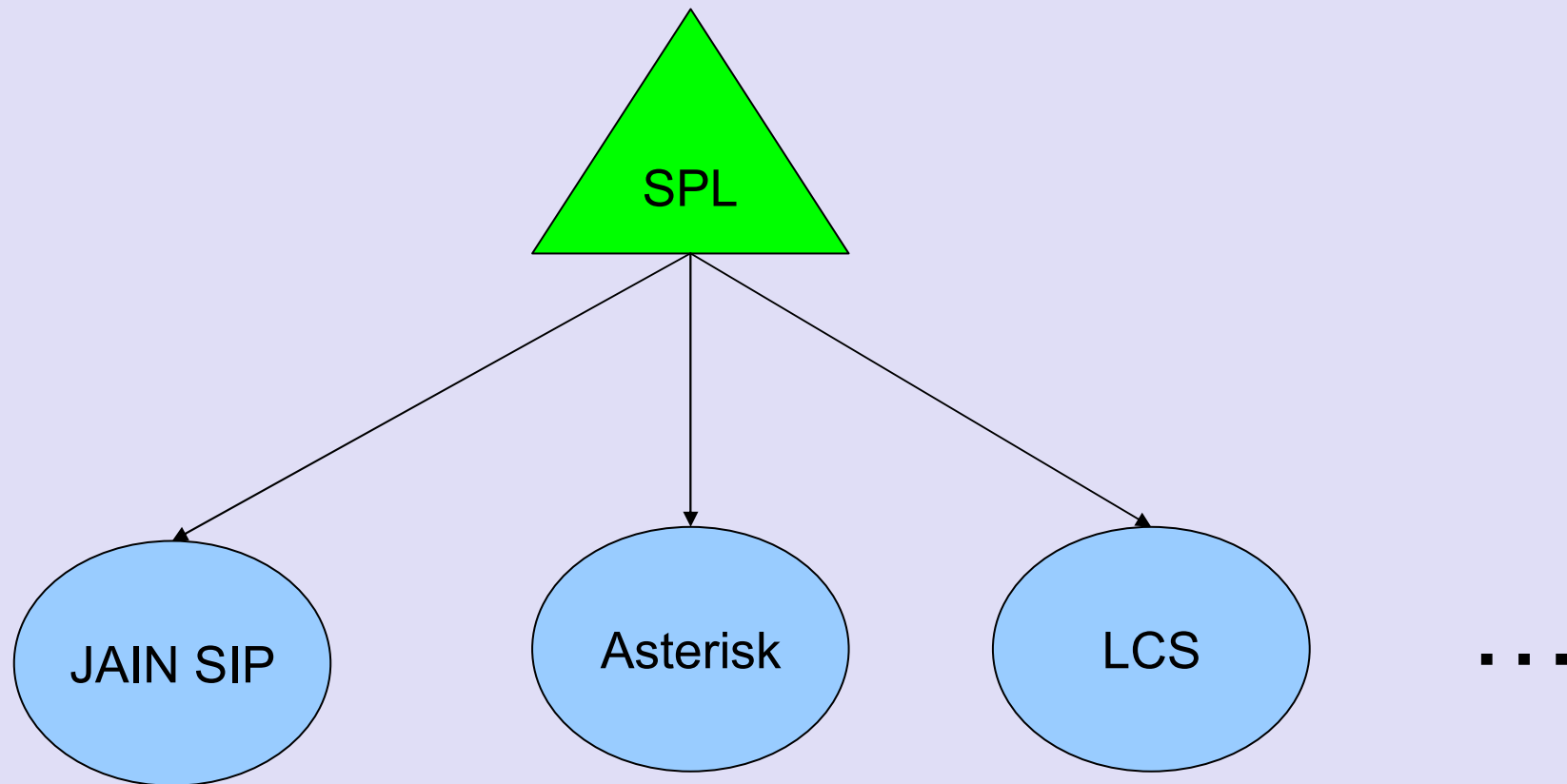
# Talk Outline

---

- ◆ Introduction to Domain Specific Languages (DSL)
- ◆ Overview of SIP
- ◆ SPL: A DSL for communication services
- ◆ **Properties of SPL**
- ◆ Summary

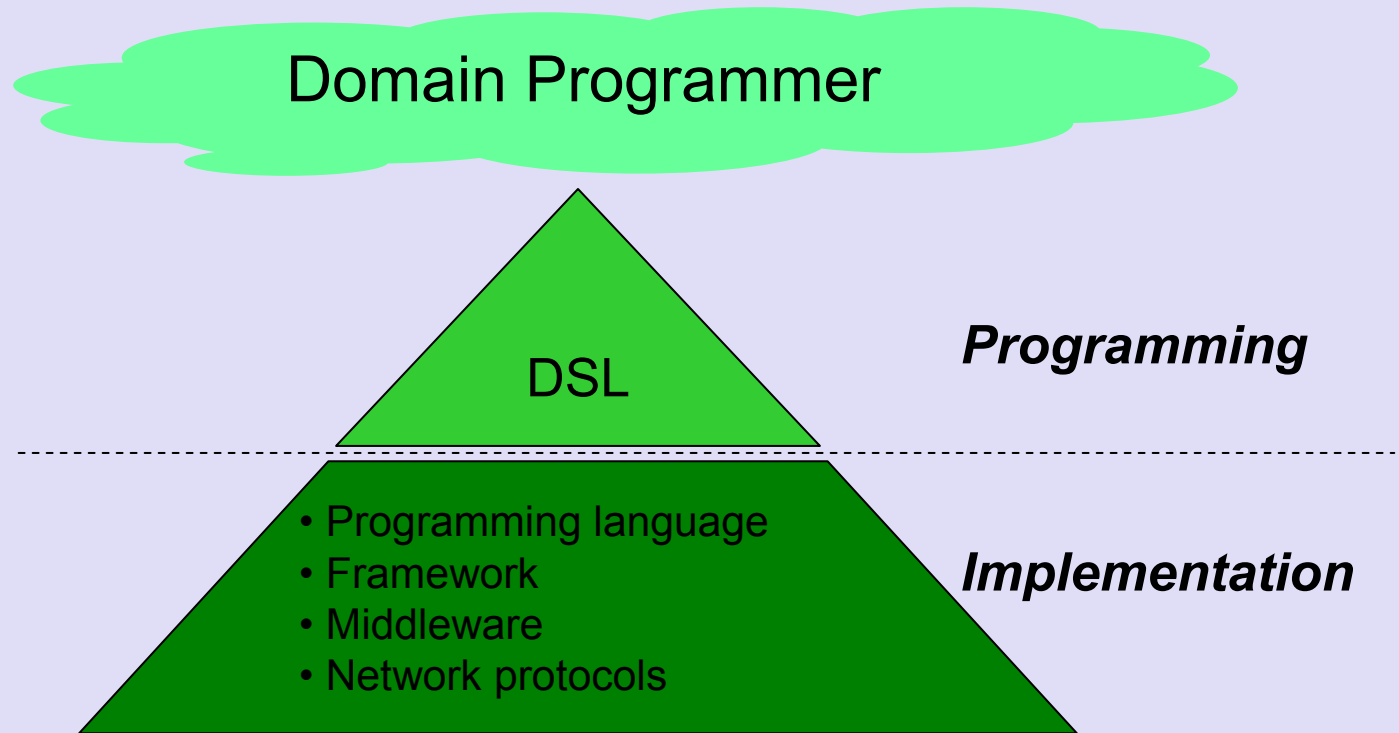
# Portability of SPL Services

---



# Abstraction Layers

---



## SPL

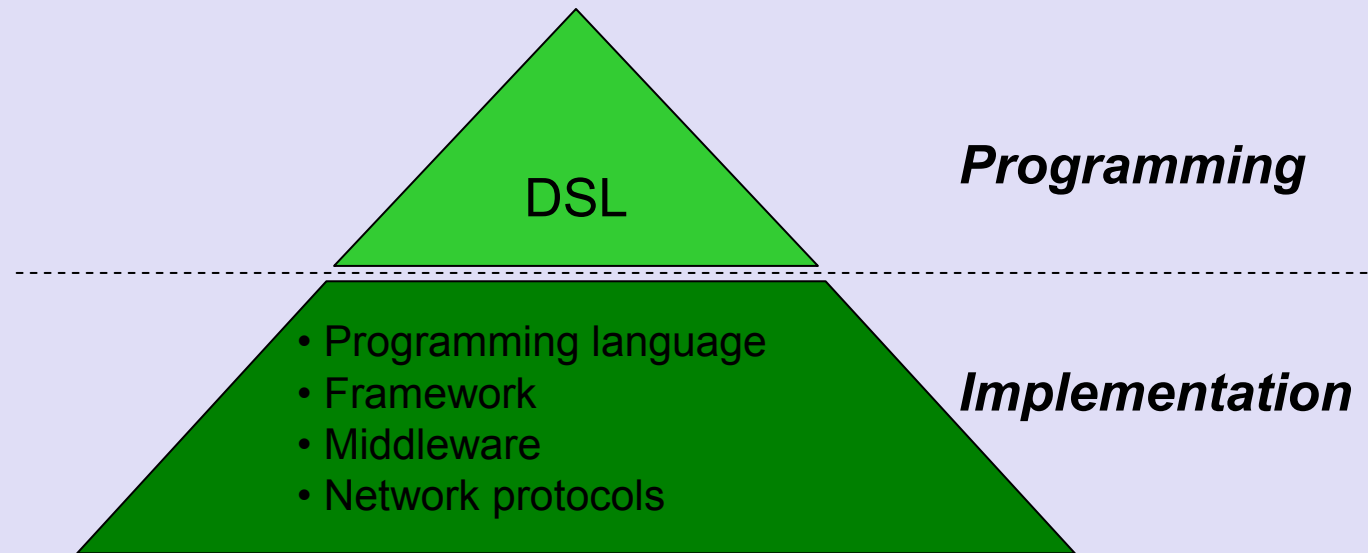
```
service example {  
  [...]  
  response incoming INVITE() {  
    response r =  
      forward 'sip:bob@phone.example.com';  
    if (r == /ERROR/CLIENT/BUSY_HERE)  
      return  
      forward 'sip:bob@voicemail.example.com';  
    else  
      if (r == /ERROR) {  
        if (FROM == 'sip:boss@example.com')  
          return forward 'tel:+19175554242';  
        return r;  
      }  
    }  
  }  
  [...]  
}
```

## ≈ JAIN SIP

```
public class Example implements SipListener {  
  [...]  
  private AddressFactory factory = getAddressFactory();  
  
  public void processRequest (RequestEvent requestEvent) {  
    Request rq_request = requestEvent.getRequest();  
    SipProvider rq_sipProvider = (SipProvider) requestEvent.getSource();  
    String method = rq_request.getMethod();  
    [...]  
    if (method.equals (Request.INVITE)) {  
      SipURI uri = factory.createSipURI ("bob", "phone.example.com");  
      rq_request.setRequestURI (uri);  
      ClientTransaction ct = rq_sipProvider.getNewClientTransaction(rq);  
      ct.sendRequest (rq_request);  
      ... }  
  
  public void processResponse (ResponseEvent responseEvent) {  
    ClientTransaction rs_ct = responseEvent.getClientTransaction();  
    if (rs_ct != null) {  
      Request rs_request = rs_ct.getRequest();  
      Response rs_response = responseEvent.getResponse();  
      SipProvider rs_sipProvider = (SipProvider) responseEvent.getSource();  
      String method = rs_request.getMethod();  
      rs_responseCode = rs_response.getStatusCode();  
      if (method.equals (Request.INVITE)) {  
        if (rs_responseCode == 486) {  
          SipURI uri = factory.createSipURI ("bob", "voicemail.example.com");  
          rs_request.setRequestURI (uri);  
          rs_sipProvider.sendRequest (rs_request);  
        } else if (rs_responseCode > 300) {  
          if (rs_request.getHeader("FROM").equals("sip:boss@example.com ")) {  
            TelURL tel = factory.createTelURL ("tel:+19175554242");  
            rs_request.setRequestURI (tel);  
            rs_sipProvider.sendRequest (rs_request);  
          } else {  
            rs_sipProvider.sendResponse (rs_response);  
          } ... }  
        }  
      }  
    }  
  }  
}
```

# Abstraction Layers

- GAP**
- ◆ Programming skills
  - ◆ Wide range of applicability
  - ◆ Programming-oriented errors
  - ◆ Programming-oriented safety properties

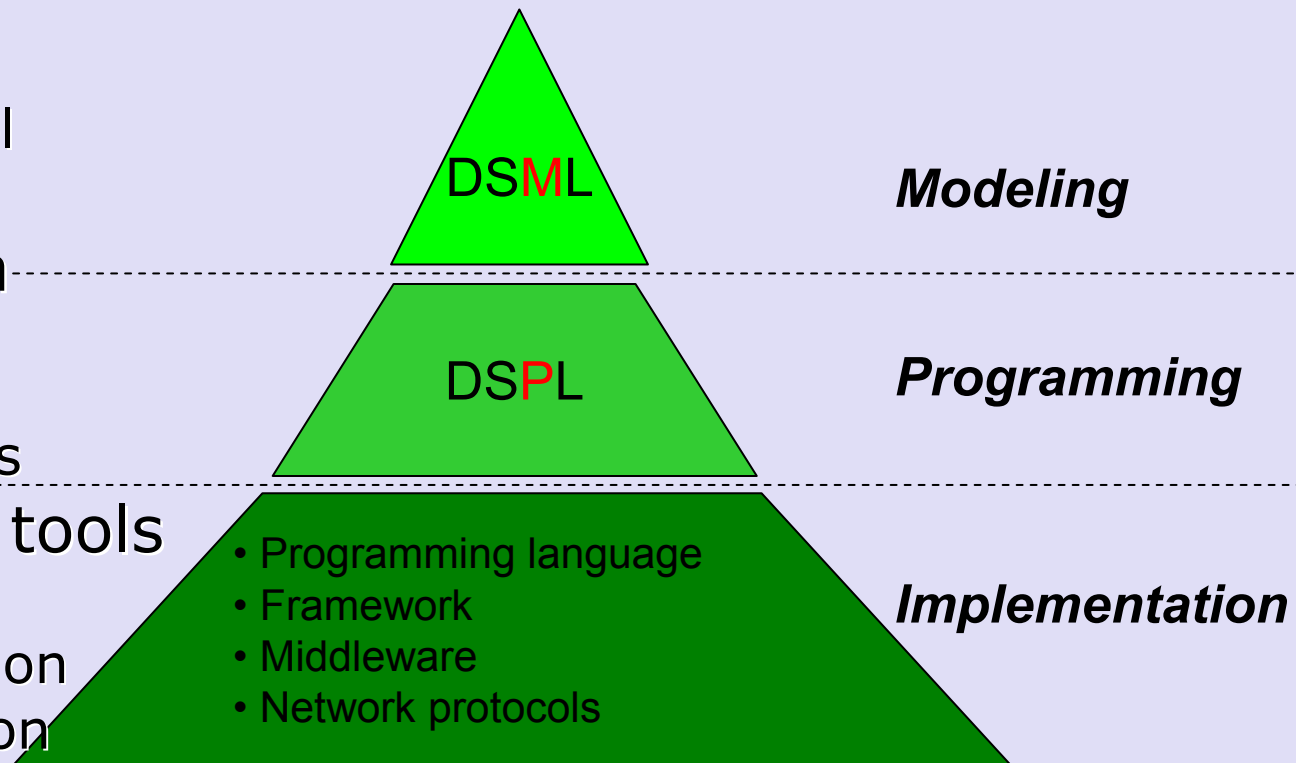




# A Layered Domain-Specific Language Approach

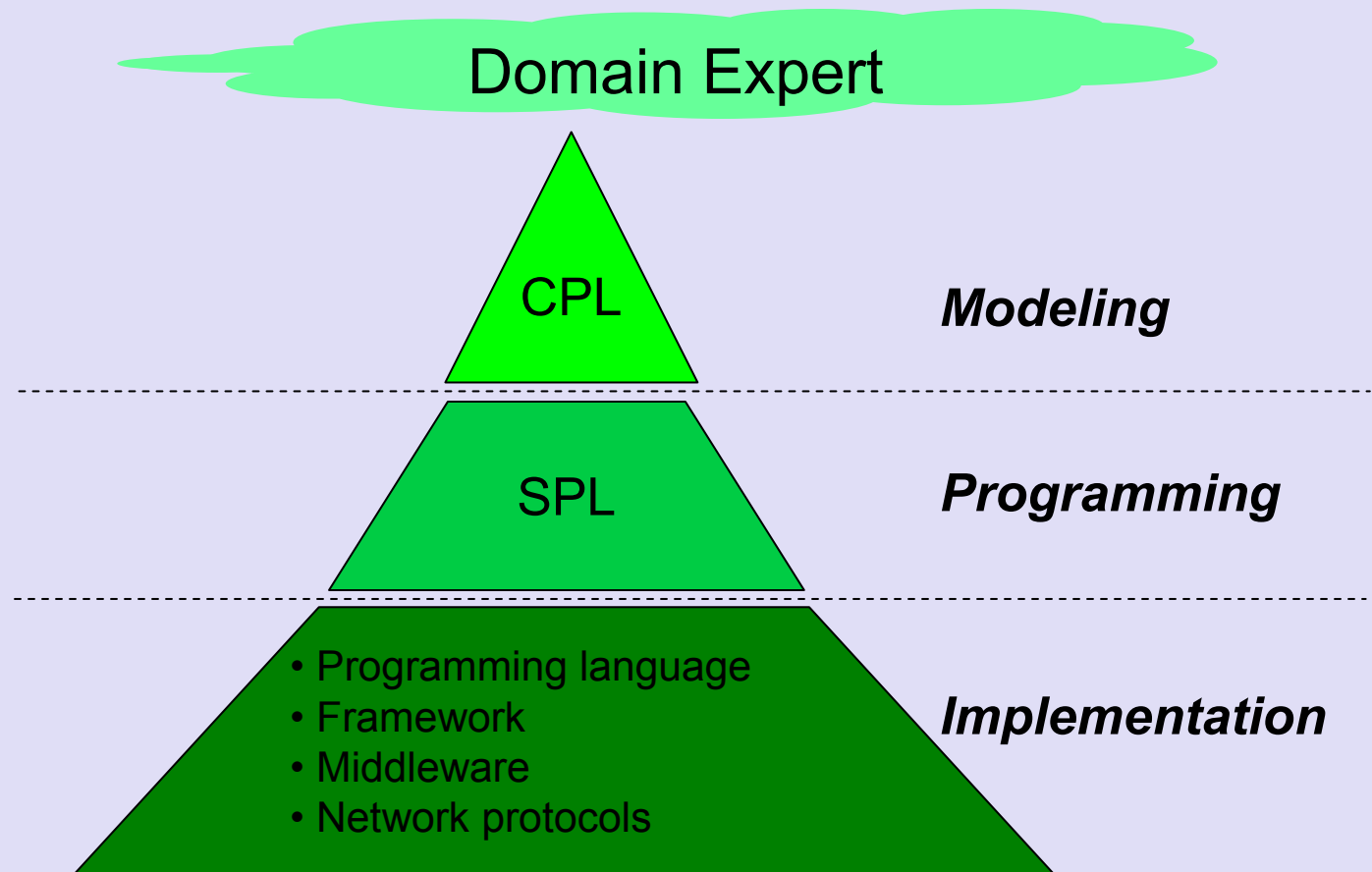
Domain Expert

- ◆ Defining a solution
  - High level
  - Simple
- ◆ Verifying a solution
  - Domain properties
- ◆ High-level tools for
  - Compilation
  - Verification

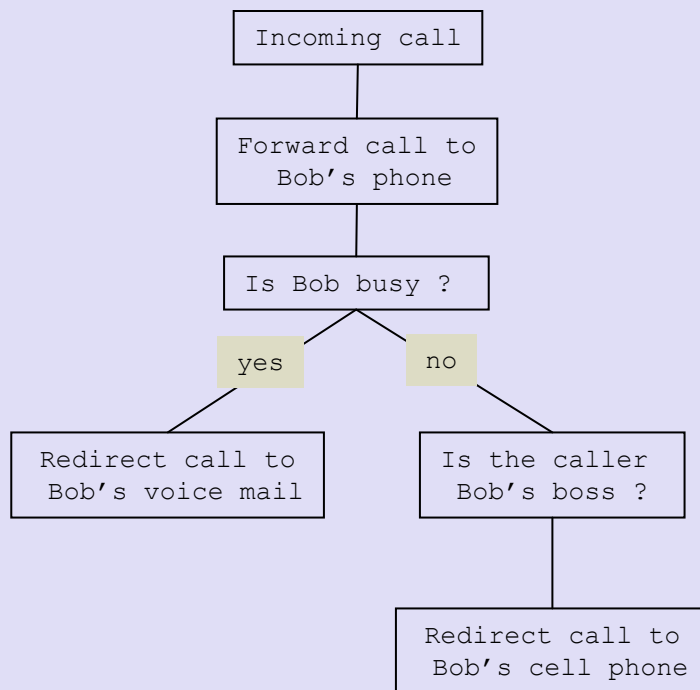


# A Layered Domain-Specific Language Approach

- ◆ Defining a solution
  - Visual
  - Activity diagram
- ◆ Verifying a solution
  - Domain properties
- ◆ Compilation
  - SPL
  - Formulas

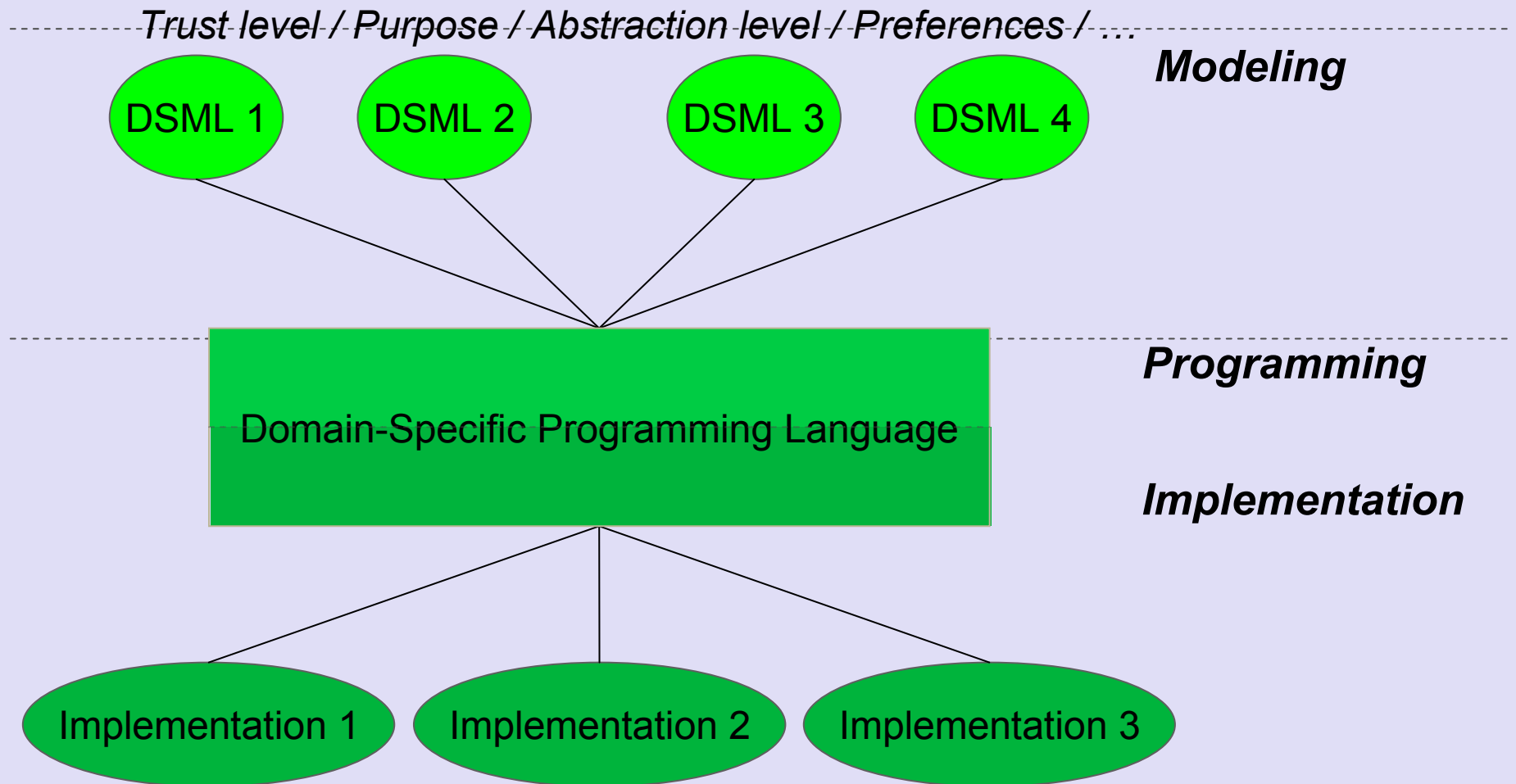


# A CPL Telephony Service



```
<?xml version="1.0" encoding="UTF-8"?>
<cpl>
  <incoming>
    <location url="sip:bob@phone.example.com">
      <proxy>
        <busy>
          <location url="sip:bob@voicemail.example.com">
            <proxy />
          </location>
        </busy>
        <otherwise>
          <address-switch field="origin">
            <address is="sip:boss@example.com">
              <location url="tel:+19175554242">
                <proxy />
              </location>
            </address>
          </address-switch>
        </otherwise>
      </proxy>
    </location>
  </incoming>
</cpl>
```

# Introducing Multiple DSMLs

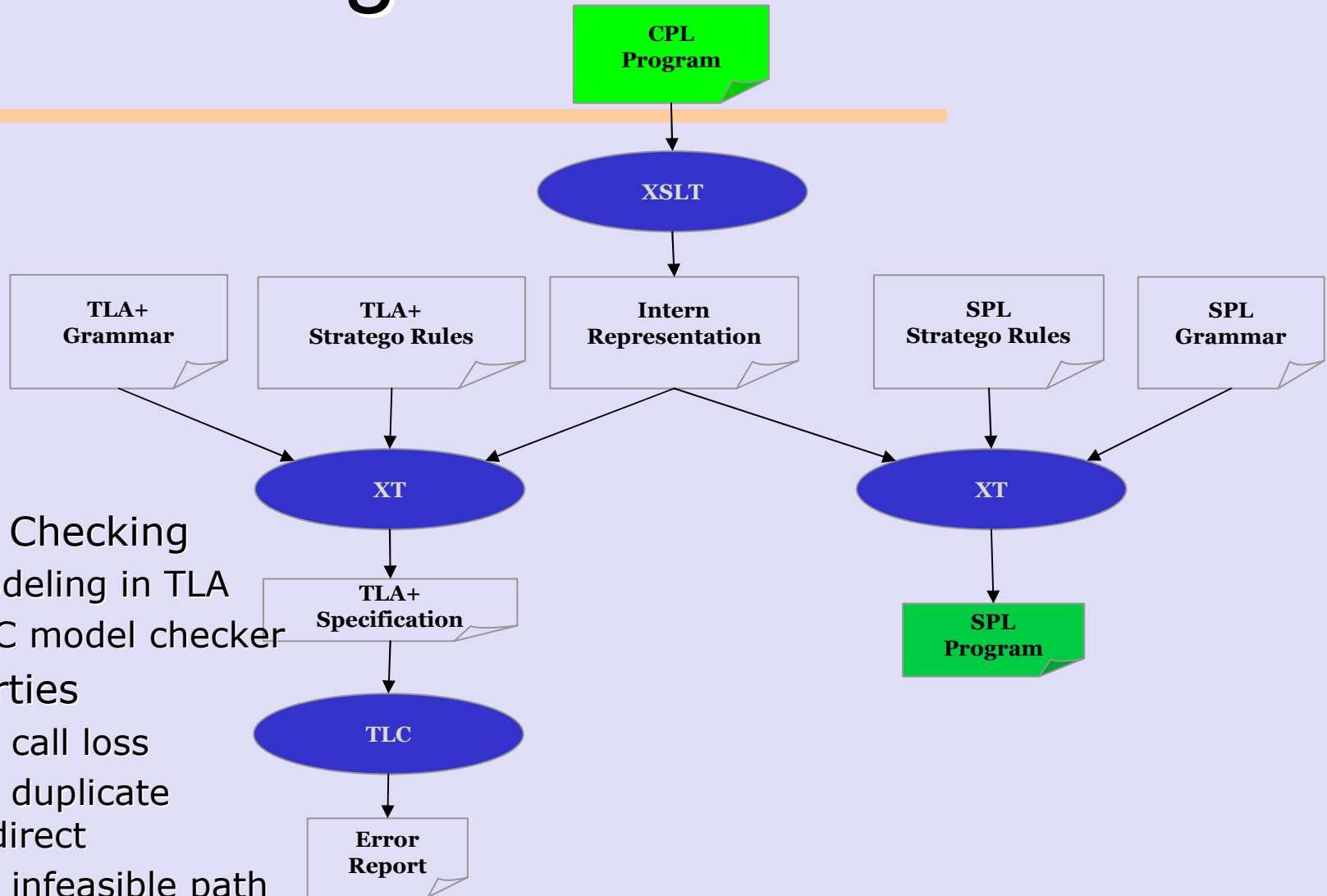


# Processing DSMLs

---

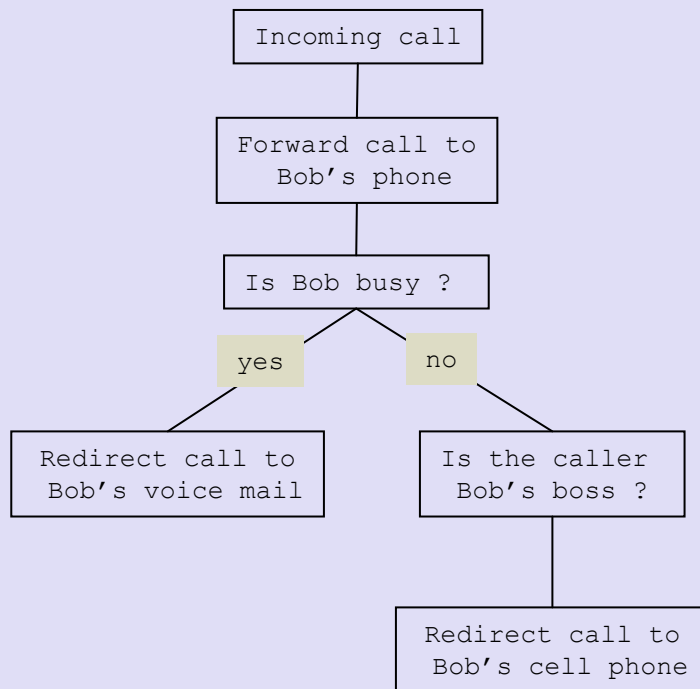
- ◆ High-level tools
- ◆ Compilation
  - DSPL Implementation
  - Formal modeling
- ◆ Verification
  - Domain-oriented properties

# Processing A DSML



- ◆ Model Checking
  - Modeling in TLA
  - TLC model checker
- ◆ Properties
  - No call loss
  - No duplicate redirect
  - No infeasible path

# From CPL to SPL



```
service example {  
  processing {  
  
    dialog {  
  
      response incoming INVITE() {  
        response r = forward 'sip:bob@phone.example.com';  
        if (r == /ERROR/CLIENT/BUSY_HERE) {  
          return forward 'sip:bob@voicemail.example.com';  
        } else if (r == /ERROR) {  
          if (FROM == 'sip:boss@example.com') {  
            return forward 'tel:+19175554242';  
          }  
          return r;  
        }  
      }  
    }  
  }  
}
```

# Domain-Oriented Properties

$AtLeastOneSigAction \triangleq currentNode = \text{"End"} \Rightarrow Len(sigActions) \neq 0$

$NoTwiceRedirectToTheSameURI \triangleq$

$\square(\forall n \in 1 .. Len(sigActions) : \forall m \in n + 1 .. Len(sigActions) :$   
 $sigActions[n] \neq \text{"Continuation"} \Rightarrow sigActions[n] \neq sigActions[m])$

$Consistency \triangleq$

$\wedge \square(\exists x \in Addresses : \forall n \in 1 .. Len(addrTest) : x \in addrTest[n])$   
 $\wedge \square(date \neq \{\})$



# Checking A Telephony Service

TLC Version 2.0 of January 16, 2006  
Model-checking

...

Finished computing initial states: 1 distinct state generated.

**Error: Invariant line 143, col 23 to line 143, col 35 of module CPL is violated.**

The behavior up to this point is:

**STATE 1:** <Initial predicate>

```

^ addrTest = □ □
^ currentNode = "Incoming"
^ sigActions = □ □
^ date = {[day ↦ "sun", dayNum ↦ 1, month ↦ "January"],
          ...
          [day ↦ "sun", dayNum ↦ 365, month ↦ "December"]}

```

**STATE 2:** <Action line 51, col 9 to line 53, col 60 of module CPL>

```

^ addrTest = □ □
^ currentNode = "WeeklyMeeting"
^ sigActions = □ □
^ date = {[day ↦ "sun", dayNum ↦ 1, month ↦ "January"],
          ...
          [day ↦ "sun", dayNum ↦ 365, month ↦ "December"]}

```

**STATE 3:** <Action line 98, col 9 to line 101, col 54 of module CPL>

```

^ addrTest = □ □
^ currentNode = "AnnualHolidays"
^ sigActions = □ □
^ date = {[day ↦ "tue", dayNum ↦ 3, month ↦ "January"],
          ...
          [day ↦ "tue", dayNum ↦ 192, month ↦ "July"],
          [day ↦ "tue", dayNum ↦ 199, month ↦ "July"],
          ...
          [day ↦ "tue", dayNum ↦ 360, month ↦ "December"]}

```

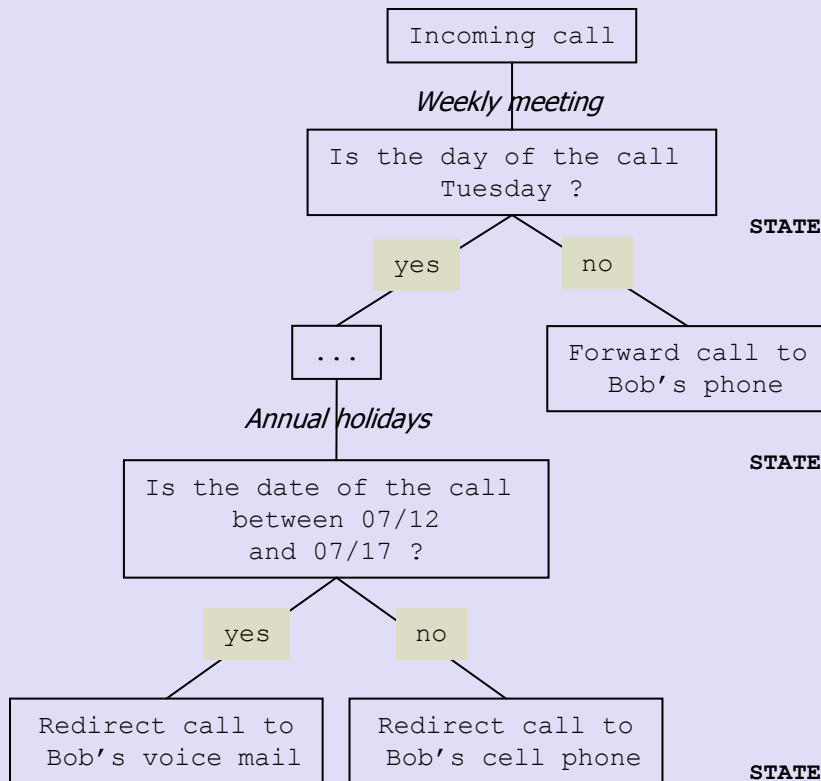
**STATE 4:** <Action line 110, col 9 to line 113, col 55 of module CPL>

```

^ addrTest = □ □
^ currentNode = "RedirectToBobVoiceMail"
^ sigActions = □ □
^ date = {}

```

5 states generated, 5 distinct states found, 2 states left on queue.  
The depth of the complete state graph search is 4.



# Summary

---

- ◆ SPL: Session Processing Language
- ◆ Contributions
  - High level abstractions and notations
    - ◆ Specific to the telephony domain
    - ◆ Ease the development process
    - ◆ Conciseness (factor of 4)
  - Performance
    - ◆ Domain-specific optimization (*e.g.*, state management)
  - Robustness
    - ◆ Critical properties guaranteed at different levels:
      - protocol, platform, and service